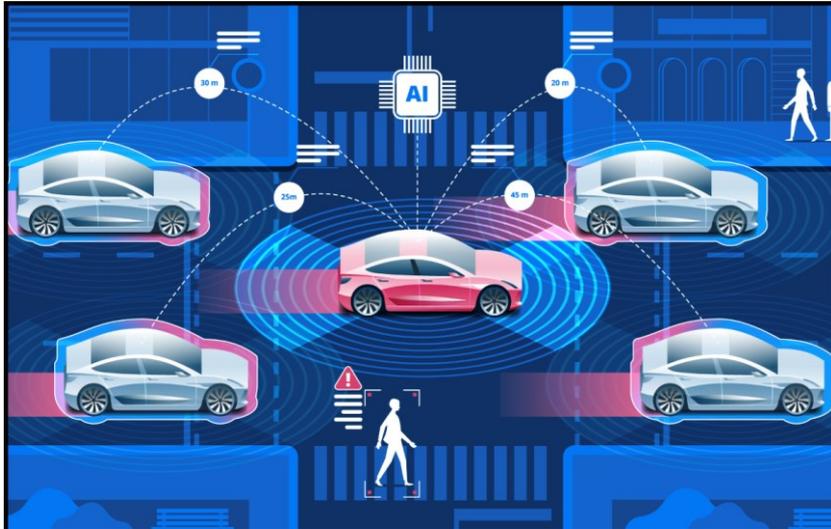


# Virtual Gang Scheduling of Parallel Real-Time Tasks

Waqar Ali<sup>\*</sup>, Rodolfo Pellizzoni<sup>+</sup> and Heechul Yun<sup>\*</sup>

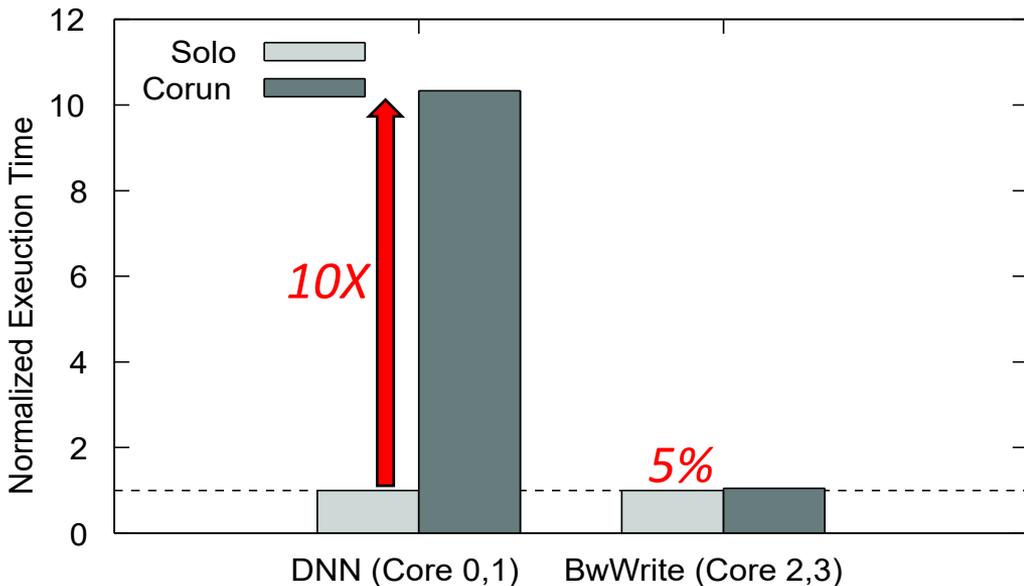
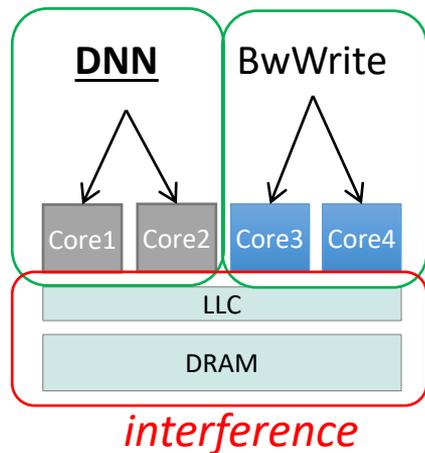
# Parallel Real-Time Applications

- Many emerging workloads are parallel real-time tasks
- Can benefit from parallelization on multicore platforms



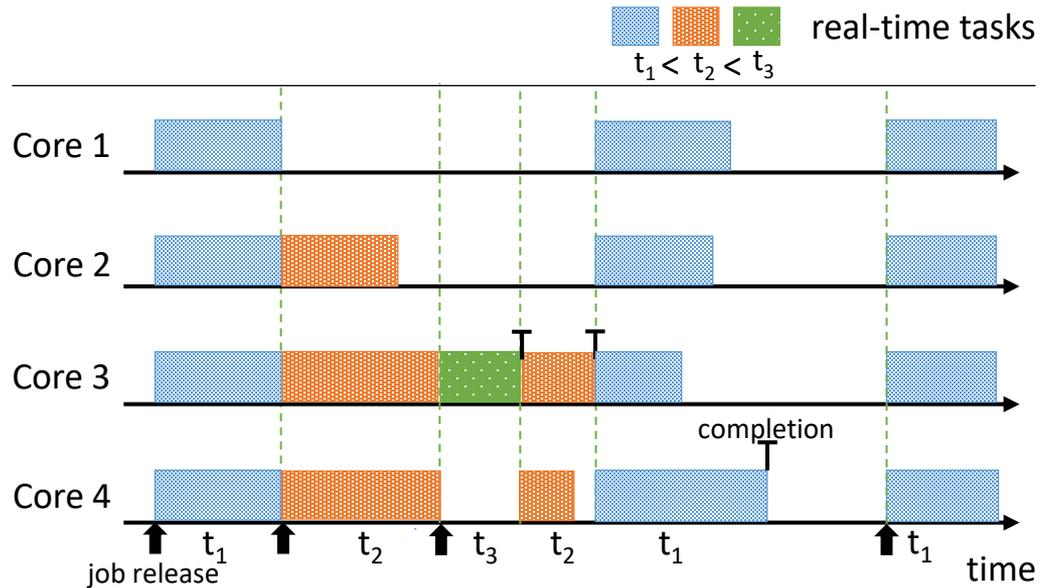
# Interference in Multicore

- Execution time of a task **depends** on co-running tasks
- Due to interference in shared hardware resources



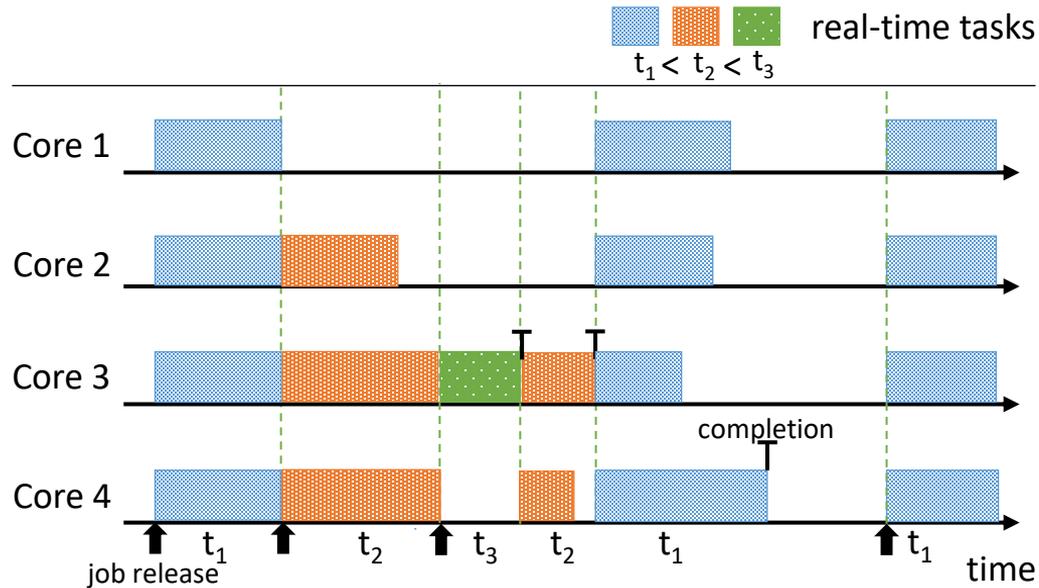
# RT-Gang

- A **gang scheduling** approach to avoid interference
- Schedule one gang task at a time



# Problems of RT-Gang

- Not all tasks are fully parallelizable
- Low utilization



# Our Approach: Virtual Gang Scheduling

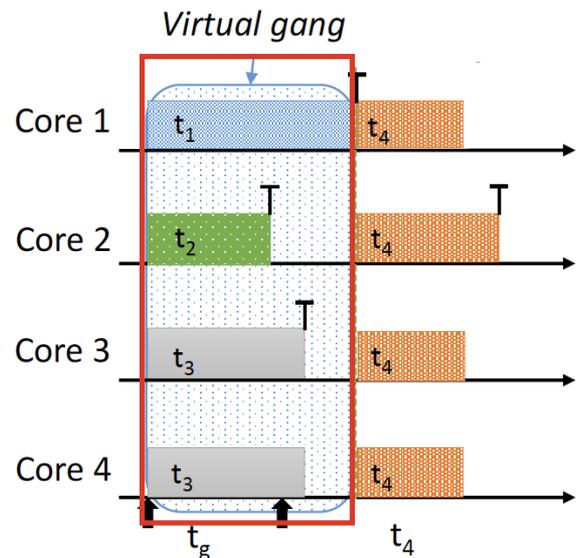
- **Virtual Gang**

- A statically paired group of real-time tasks
- Scheduled as a single gang task

- **Benefits**

- Can achieve higher utilization
- Co-running tasks never change
- Can tightly bound interference

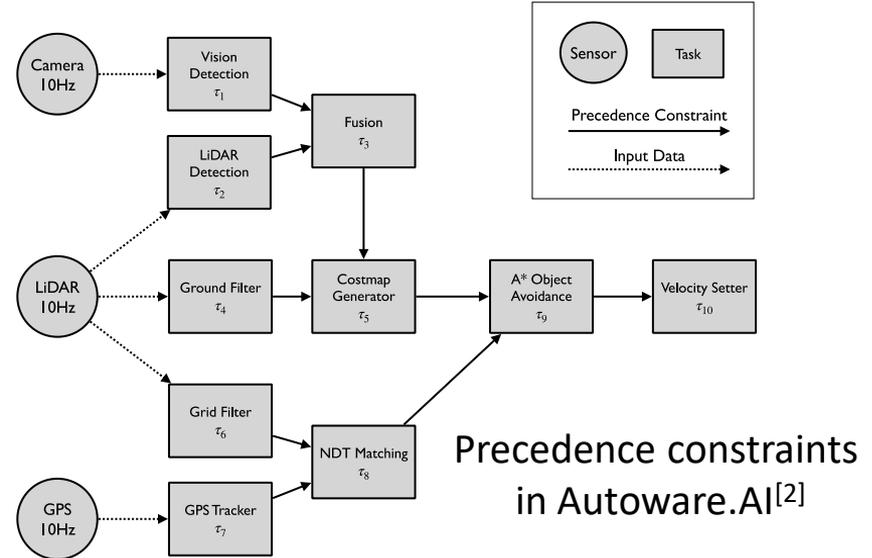
- **Q: How to form the virtual gangs?**



Real-Time Tasks:  $t_1, t_2, t_3, t_4$   
Virtual Gang ( $t_g$ ):  $t_1 + t_2 + t_3$   
 $\text{prio}(t_g) > \text{prio}(t_4)$

# Considerations

- **Interference**
  - Co-running tasks influence execution timing
- **Precedence constraints**
  - Common in real applications
  - Limit feasible task pairings

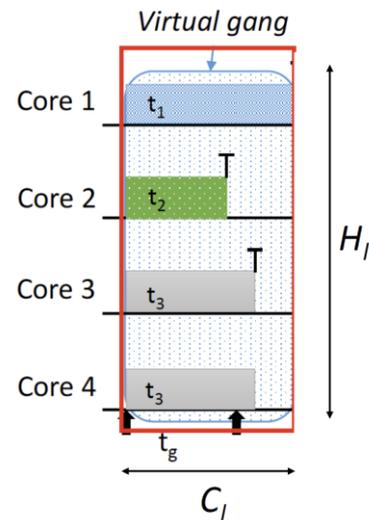


# System Model

- **A multicore platform with  $m$  unit-speed cores**
- **$n$  periodic, rigid real-time gang tasks  $\Gamma = \{\tau_1, \tau_2, \dots, \tau_n\}$**
- **A real-time gang task:  $\tau_i = (c_i, h_i, r_i, T_i)$** 
  - $c_i$  : WCET in isolation
  - $h_i$  : #of cores required to run ( $1 \leq h_i \leq m$ )
  - $r_i$  : Resource demand  $r_i \in [0,1]$
  - $T_i$  : Period of  $\tau_i$
- **Precedence constraints described by a set of DAGs**

# System Model

- **A candidate-set:**  $\Delta_T = \{\forall \tau_i \in \Gamma \mid T_i = T\}$ 
  - Tasks that share a common period  $T$
  - A small number of candidate-sets exists in  $\Gamma$
- **A virtual gang:**  $w_l = (C_l, H_l, R_l, T_l)$ 
  - A subset of tasks within a candidate set  $\Delta_T$
  - $C_l$ : WCET of the virtual gang
  - $H_l$ : collective #core demand ( $1 \leq H_l \leq m$ )
  - $R_l$ : collective resource demand ( $R_l \geq 0$ )
  - $T_l$ : common shared period



# Interference Model

- **A virtual gang  $w_l$ 's WCET:**  $C_l = \max_{\forall \tau_k \in w_l} \{c_k\} \times \max(R_l, 1)$ 
  - $R_l < 1$ : suffers no interference until the resource is over-utilized
  - $R_l > 1$ : applies a *linear scaling*

# Virtual Gang Formation Problem

- 
- For a given candidate set of  $N$  real-time tasks with the period  $T$
  - Form a set of virtual gang tasks that minimize completion time
  - Subject to a given set of precedence constraints
-

# Optimal SMT Algorithm

- **Based on Satisfiability Modulo Theories (SMT)**
  - **Step-1**: Identify the parameters of the optimization problem and the *optimization variable*
  - **Step-2**: Write the constraints that must be satisfied by the parameters in a *feasible* solution
  
- **See the paper for the details.**

# Greedy Heuristic Algorithm

- High-level idea
  - Group tasks with similar WCET values while minimizing combined resource utilization
  - higher core utilization, less slowdown
- Algorithm
  - Sort tasks by WCET
  - Pick the longest
  - Find feasible co-runners
  - Rank the feasible co-runners considering resource utilization
  - Form a virtual gang and repeat
- Not optimal but fast and effective

---

**Algorithm 1:** Virtual Gang Formation Heuristic

---

```
1 Input: Candidate Set ( $\Delta_T$ ), Number of Cores ( $m$ )
2 Output: Taskset comprising virtual gangs
3 function gang_formation( $\Delta_T, m$ )
4    $pq = \text{sort\_tasks\_by\_wcet}(\Delta_T)$ 
5    $\text{virtualGangs} = ()$ 
6   while not_empty( $pq$ ) do
7      $\tau_i = pq.\text{pop}()$ 
8      $f_i = \text{family}(\tau_i)$ 
9      $\text{partners} = ()$ 
10    for  $\tau_j \in pq$  do
11      if  $\tau_i.h + \tau_j.h \leq m \wedge \tau_j \notin f_i$  then
12         $\text{partners} \leftarrow \text{partners} \cup \{\tau_j\}$ 
13     $pq_i = \text{score\_partners}(\text{partners})$ 
14    while not_empty( $pq_i$ ) do
15       $\tau_p = pq_i.\text{pop}()$ 
16       $\tau_i = \text{merge}(\tau_i, \tau_p)$ 
17       $pq.\text{remove}(\tau_p)$ 
18       $\text{update\_partners}(\tau_i, pq_i)$ 
19     $\text{virtualGangs} \leftarrow \text{virtualGangs} \cup \{\tau_i\}$ 
20 return  $\text{virtualGangs}$ 
```

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# Response Time Analysis

- For a taskset of virtual gangs  $\{w_1, w_2, \dots, w_l\}$  and under the rate-monotonic priority assignment, the response time  $\mathbb{R}_i$  of gang  $w_i$  can be iteratively computed using:

$$\mathbb{R}_i^{k+1} = \bar{C}_i + \sum_{\forall w_j \in hp(w_i)} \left\lceil \frac{\mathbb{R}_i^k}{T_j} \right\rceil C_j$$

- $\bar{C}_i$  is the WCET of  $w_i$  + the WCET of all  $w_j$  (with the same period  $T$  as  $w_i$ ) that come before  $w_i$  in the linear execution order

First work which enables schedulability analysis of real-time gang tasks with precedence constraints!

# Evaluation

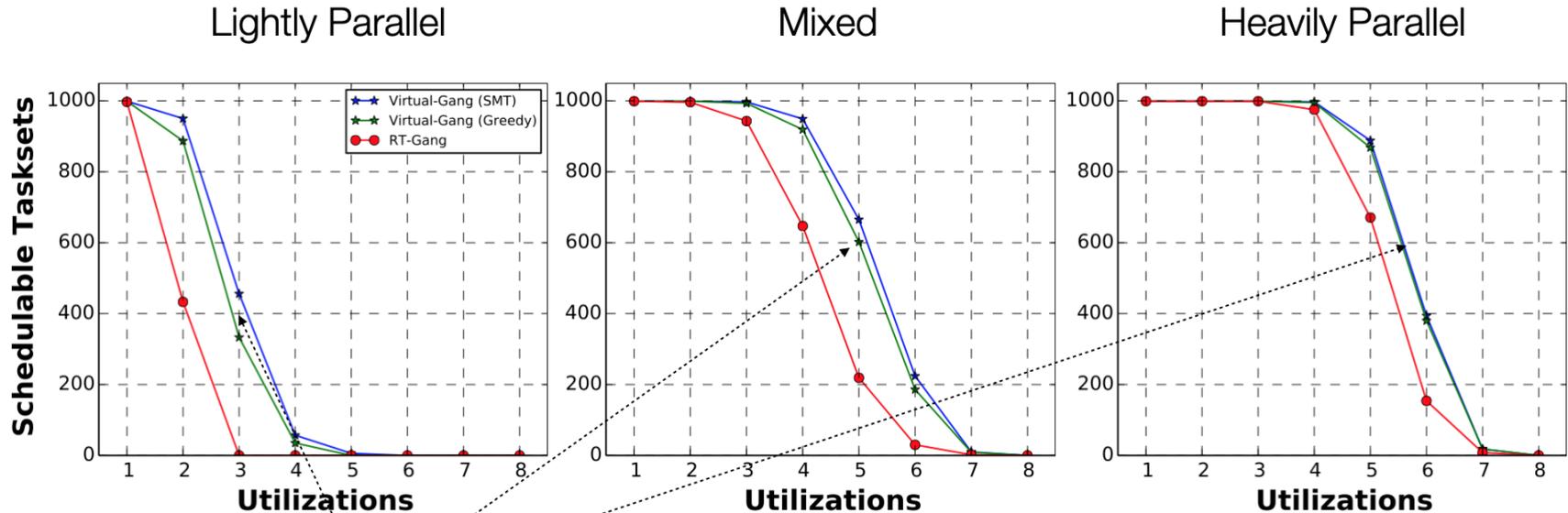
- **Simulation study with randomly generated synthetic tasksets**
  - **Lightly parallel** ( $h=[1, \lceil 0.3 \times m \rceil]$ )
  - **Mixed** ( $h=[1, m]$ )
  - **Heavily parallel** ( $h=[\lceil 0.3 \times m \rceil, m]$ )
  - **See the paper for other parameters (period, tasks, resource demand, etc.)**
  
- **Compared scheduling schemes**
  - **RT-Gang: one gang task at a time** <sup>[1]</sup>
  - **Virtual Gang (SMT/Greedy): our approach**
  - **Gang-FTP: fixed-priority gang scheduling** <sup>[3]</sup>
  - **Threaded: vanilla Linux (global FP scheduling)** <sup>[4]</sup>

[1] W. Ali and H. Yun, "RT-Gang: Real-Time Gang Scheduling Framework for Safety-Critical Systems," In RTAS (2019)

[3] S. Wasly and R. Pellizzoni, "Bundled Scheduling of Parallel Real-Time Tasks," In RTAS (2019)

[4] J. Fonseca et al., "Improved Response Time Analysis of Sporadic DAG Tasks for Global FP Scheduling," In RTNS (2017)

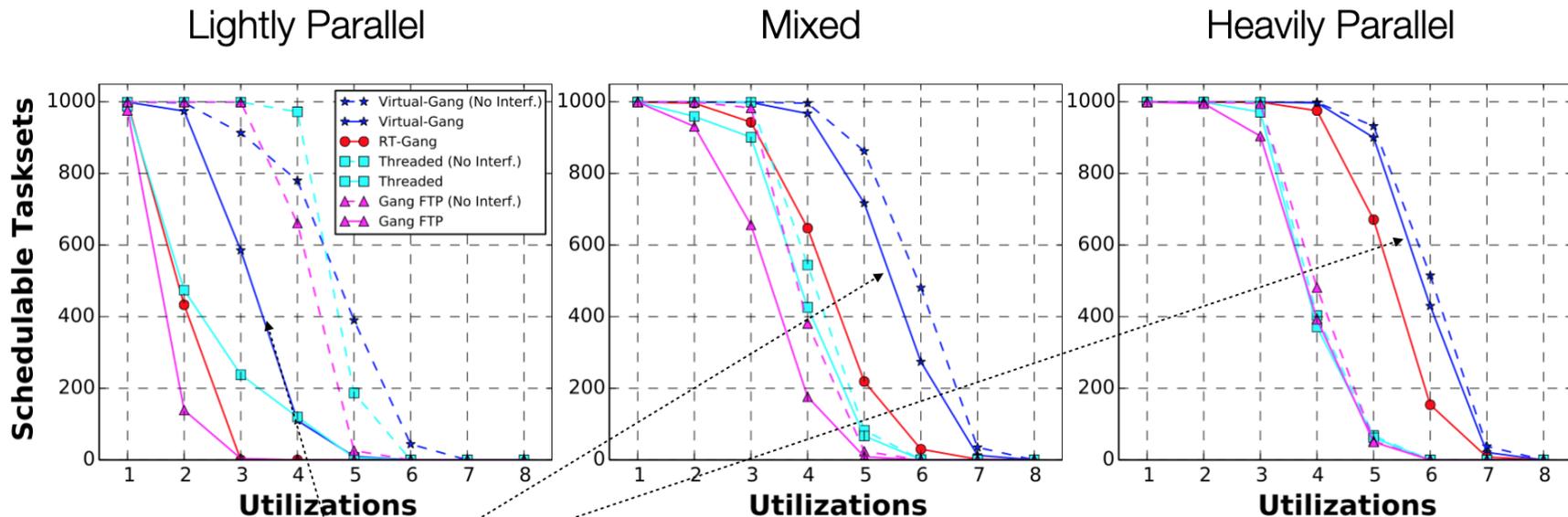
# Results (with Precedence Constraints)



▶ Heuristic performs very close to the optimal algorithm!

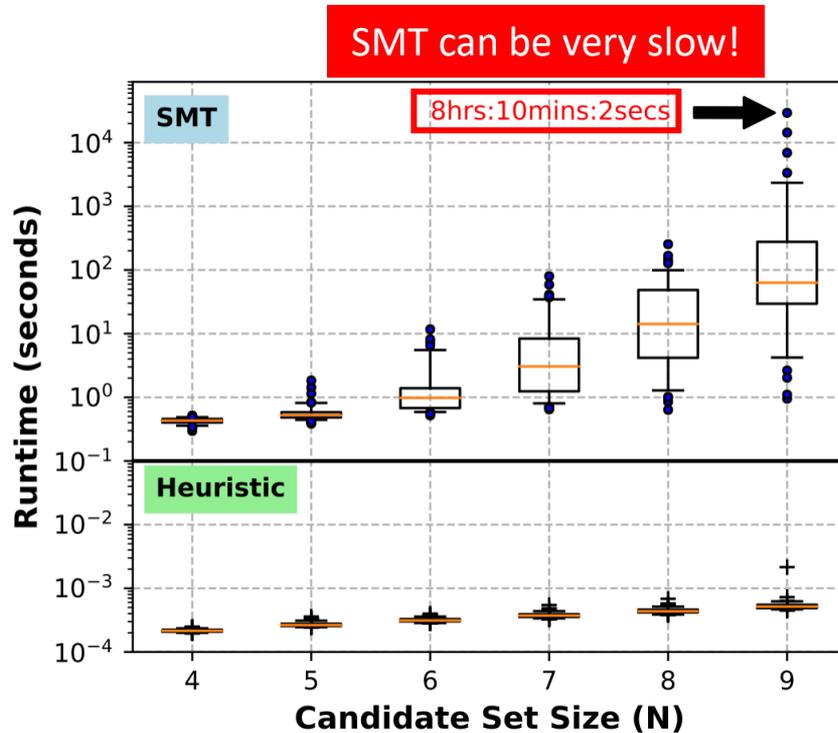
▶ Schedulability is highest for virtual-gang (SMT)

# Results (w/o Precedence Constraints)



- ▶ When interference is considered:
  - ✓ Virtual Gang scheme outperforms every other analysis for all taskset types

# Runtime of Gang Formation Algorithms



Even for  $N = 50$ , the heuristic gives a solution in less than a second!

# Summary

- **Proposed a virtual-gang scheduling scheme**
  - Achieves high utilization
  - Considers interference and precedence constraints
  - Enables effective and tight timing analysis
- **Future work**
  - Virtual gang formation on heterogeneous platforms (w/ accelerators)
  - Empirical evaluation with real-world workloads

<https://github.com/CSL-KU/VirtualGang-Simulator>