

# Resource Management and Differentiated Services

Lixia Zhang

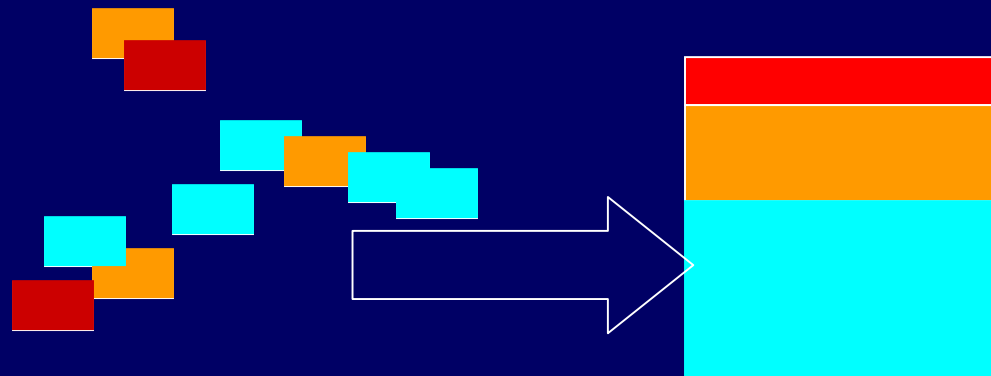
UCLA Computer Science Department

with input from many others

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# Network QoS support seen from 10,000 feet:

- 1 Define packet treatments at switches/routers
- 2 Control the amount of resources allocated to each treatment class
- 3 Sort packets into classes



# What this talk is about

- ◆ How to provide scalable, robust, and manageable resource management
  - ▶ Ongoing research/development effort
  - ▶ incorporating others visions/ideas

## What this talk is *not* about

- ◆ how many different traffic classes needed, or how to set the TOS field value
  - ▶ the charters of IETF intserv & diff-serv Working Groups

# What is pushing diff-serv effort

- ◆ A market need since yesterday: simple mechanisms that can be *quickly* and *incrementally* deployed to provide differentiated services
  - ▶ no one wants to sell “bad” services
  - ▶ everyone wants to sell varying levels of “good” services
- ◆ Various doubts on feasibility of intserv framework
  - ▶ Complexity?
  - ▶ Scalability?
  - ▶ Quick deploy-ability?

# How and how much do diff-serv & intserv differ?

- ◆ Where to start:
  - ▶ defining end-to-end services, vs.
  - ▶ defining packet treatments at individual components

A sidenote:

- ▶ IP started by defining hop-by-hop forwarding, rather than end-to-end delivery service

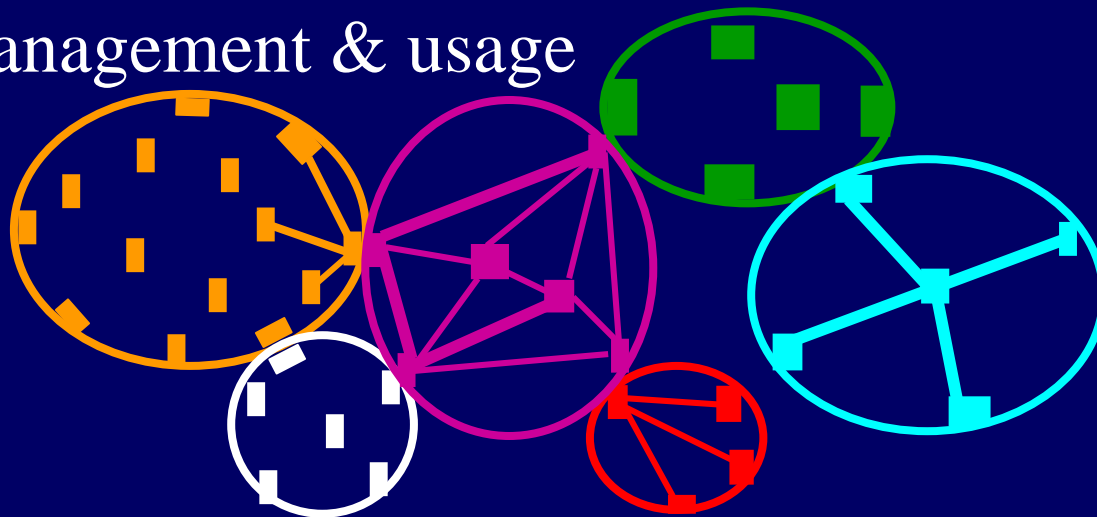
# How and how much diff-serv & intserv differ (II)

- ◆ How to control the amount of resources allocated
  - ▶ end-to-end QoS support requires end-to-end signaling
  - ▶ per-hop treatment can work with either static configuration or dynamic signaling
- ◆ How to sort packets into classes
  - ▶ old RSVP way:
    - identify individual flows
    - map packets of each flow to proper traffic classes
  - ▶ diff-serv: use TOS field as class ID
    - pre-classified somewhere

# Network resource management

Emerging model:

- ◆ interconnects of administrative domains
- ◆ a priori bilateral agreement between neighboring domains
- ◆ each domain responsible for its internal resource management & usage



# An analogy to global routing

- ◆ Hierarchical
  - ▶ needed for **scaling**
  - ▶ needed for **administrative control**
  - ▶ different granularity at different levels
- ◆ routes are *pre-computed* (or *pre-configured*)
- ◆ concatenation of hop-by-hop forwarding provides end-to-end data delivery
- ◆ routes dynamically adjustable
  - ▶ adapt to topology/policy changes



# A proposed picture for Scalable QoS support

- ◆ Two-tier resource management
  - ▶ inter-administrative domains
  - ▶ intra-administrative domains
- ◆ Inter-domain: pre-negotiated neighboring relation
  - ▶ infeasible to set up business relation upon every new flow in real-time?
- ◆ concatenation of bilateral agreement leads to end-to-end QoS delivery paths
- ◆ amount of resources adjustable
  - ▶ adapt to demand/policy/topology changes

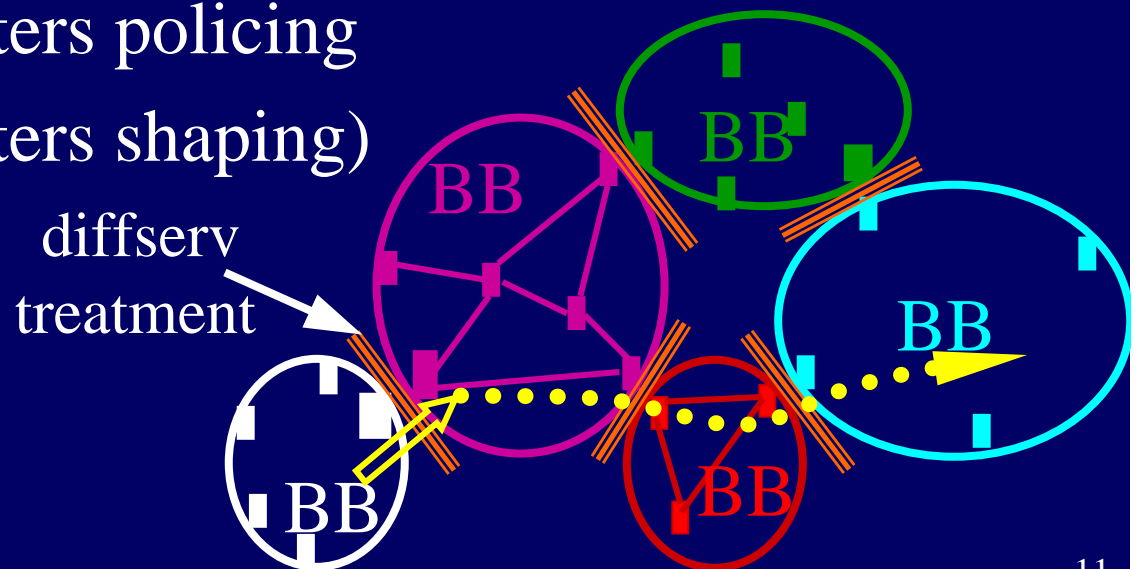
# Resource manager: Bandwidth Broker (BB)

- ◆ A logical entity residing in each administrative domain
  - ▶ Managing internal demands & resources according to the policy database (who can do what when)
  - ▶ setting up & maintaining bilateral agreement with neighbor domains
    - bookkeeping how much traffic entering which border router & going out which border router
- ◆ today's BB: network administrators & operators
  - ▶ would like to automate over time

“A Two-bit differentiated services architecture for the Internet”  
Nichols, Jacobson, Zhang  
draft-nichols--diff-arch-00.txt, November 1997

# An overall picture

- ◆ “Keep complexity at edges, leave the core simple”
  - ▶ peripheral domains may manage internal traffic and resources in any way they wish
  - ▶ border-crossing packets carrying right TOS value and treated diffserv way
- ◆ ingress border routers policing
- ◆ (egress border routers shaping)



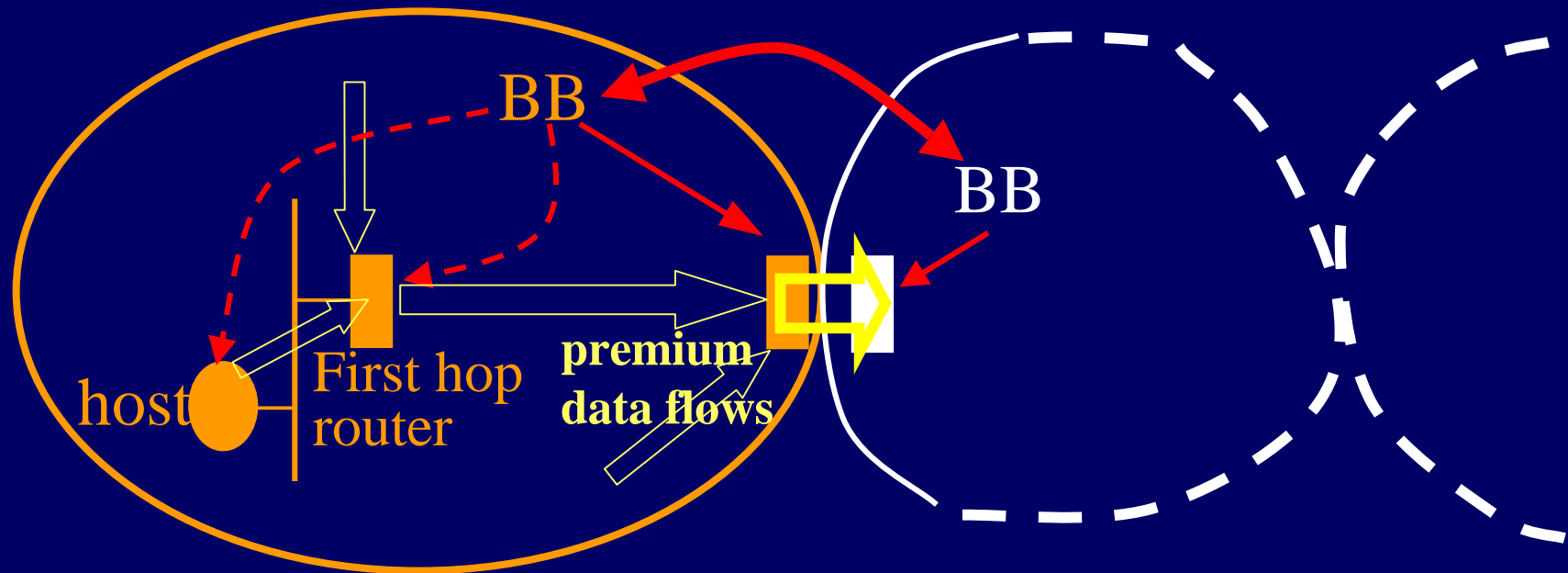
# Some of the questions

- 1 How does a leaf domain BB know the total local demands for each egress border router?
- 2 How does a transit domain BB map its inter-domain commitment to internal resource allocation?
- 3 How much (& what) state must BB keep?
- 4 How much (& what) state must a router keep?
  - ▶ Router in leaf domains
  - ▶ Router in core networks

# Choices for implementation

- ◆ adequate provisioning
  - ▶ eliminate Questions 2, 3, & 4
- ◆ manual configuration
  - ▶ not that different from static routing
- ◆ using some setup protocols
  - ▶ inter-domain: BB-to-BB
  - ▶ Intra-domain: RSVP as a ready candidate

# An example of provisioning in a local domain

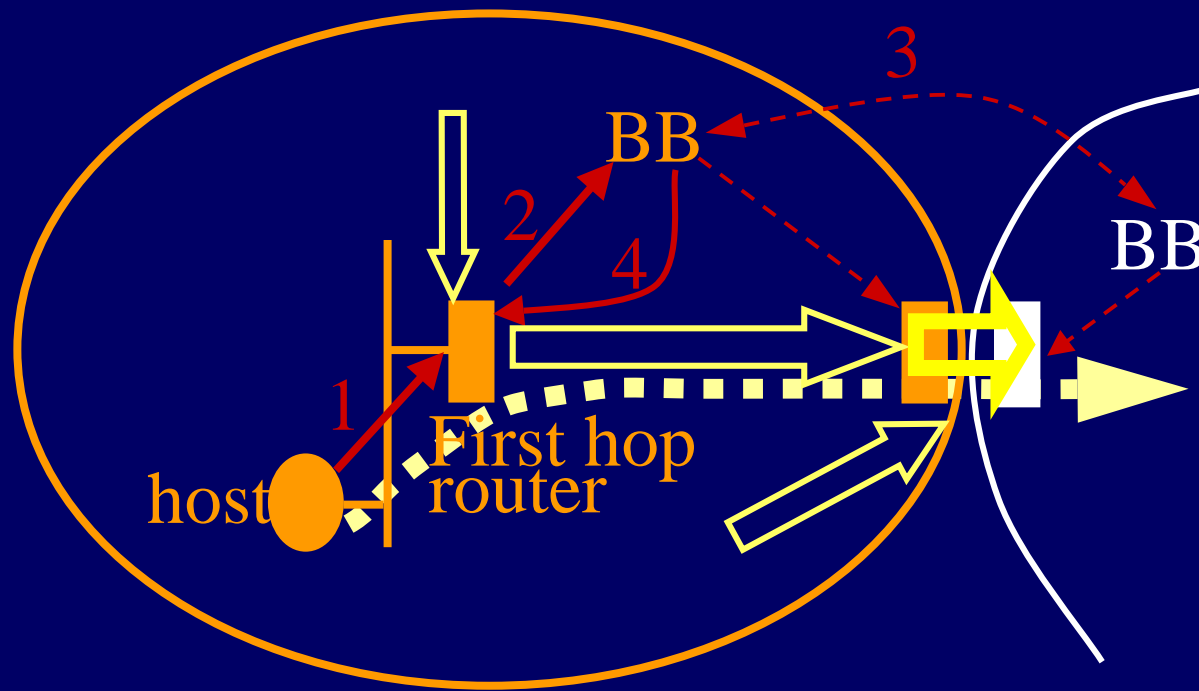


**BB** is assumed to have adequate knowledge about internal demand; may readjust the allocation over time.

**BB** & **BB** instruct their edge devices how to shape/police.

-- ➔ Indicating additional configurations (shaping/policing) if the domain cannot solely rely on provisioning.

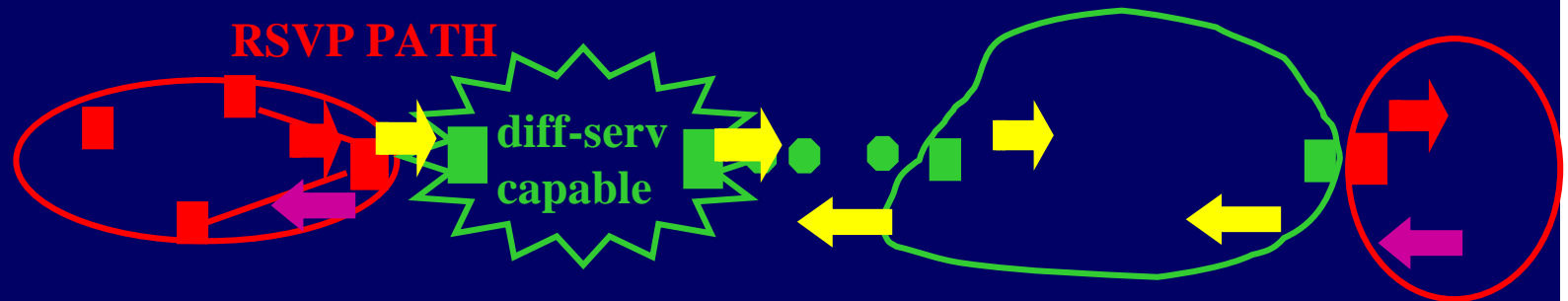
# An example of using RSVP in a local domain



**BB** may pre-reserve adequate bandwidth with BB to avoid readjusting the inter-domain allocation everytime (the actions indicated by the dashed lines)

# “Tunnel” RSVP messages between leaf domains

- ◆ Why “tunnel through”: do *not* want intermediate routers to see/act on end-to-end RSVP messages
- ◆ One way of doing it



- ◆ drawback
  - ▶ assuming both ends using RSVP internally
  - ▶ intra-domain signaling msgs crossing boundaries

“A Framework for End-to-End QoS Combining RSVP/Intserv and Differentiated Services”  
draft-bernet-intdiff-00.txt



# Intra-transit domain implementation

## Choices of implementation

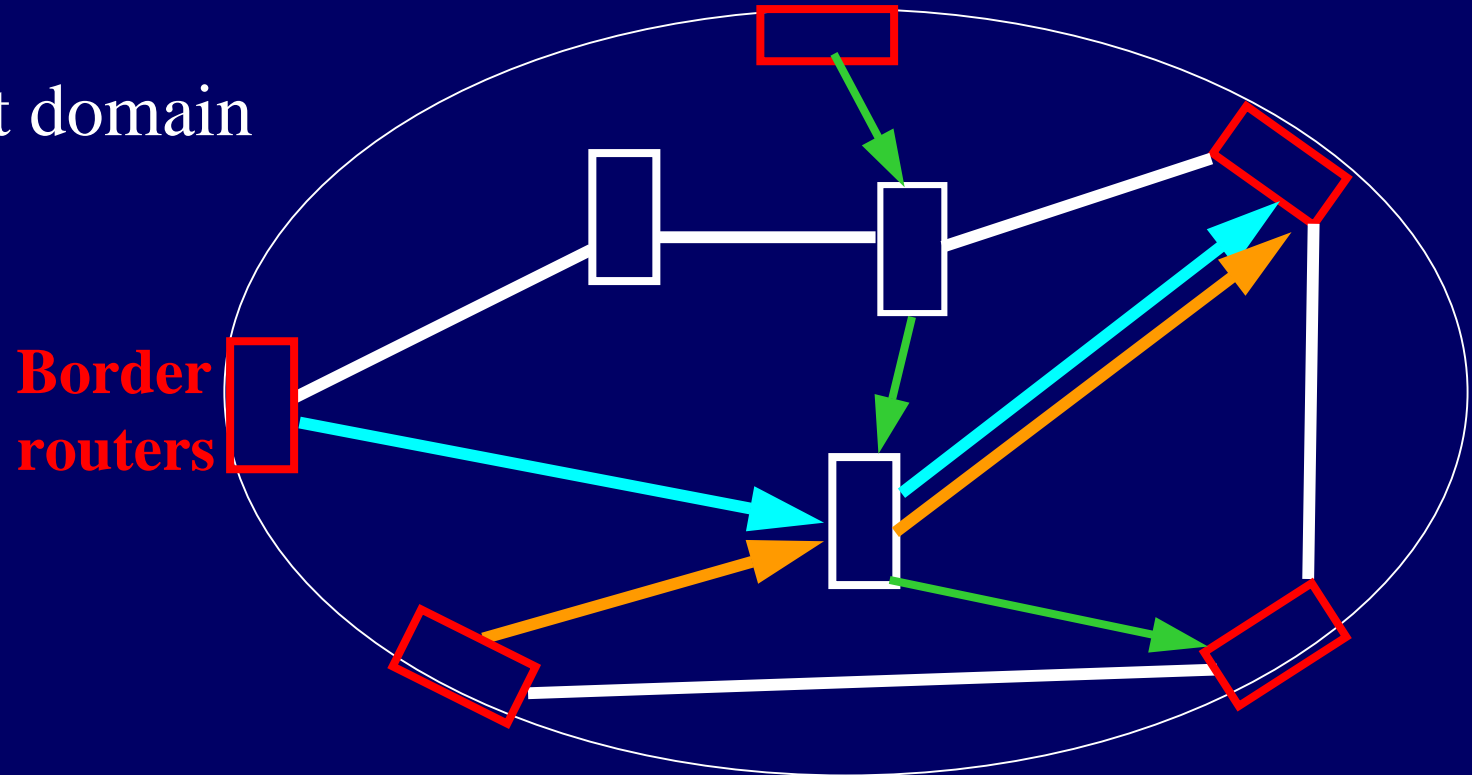
- ◆ provisioning
- ◆ manual configuration, or SNMP
- ◆ use an automatic setup protocol, such as RSVP

## How to use RSVP in core networks

- ◆ border routers behave as sources & destinations for ingress and egress traffic
  - ▶ similar idea discussed in PASTE draft

# Here is a picture

A transit domain



- Set up an RSVP session for each ingress flow
  - RSVP msgs tell each router along the way how much to reserve
  - Routers classify packets by TOS field
- Reservations follow routing changes automatically

# Some of FAQ's

- ◆ Is this sender or receiver driven?
  - ▶ At BB level: yes to all, sender domain BB, receiver domain BB, possibly 3rd party BB
- ◆ How to handle allocation for multicast traffic?
  - ▶ See above
    - details being worked out
- ◆ relation between the two levels of resource control?
  - ▶ Inter-domain (BB) level:
    - independent from whether one does anything internally
  - ▶ intra-domain: local decision

## Summary: One possible picture for **diff-serv resource management**

- ◆ two-level hierarchy
  - ▶ inter-domain management
    - currently human
    - automate over time; BB as one proposal
  - ▶ intra-domain management: multiple possible choices
    - provisioning, manual-configuration, SNMP, RSVP
- ◆ packet classification
  - ▶ cross-domain traffic: classified by bits in TOS field
  - ▶ leaf domain: one's own choice
- ◆ Work underway for a prototype implementation