# <u>The Binomial Multi-Section</u>

# Transformer

Recall that a **multi-section matching network** can be described using the theory of small reflections as:

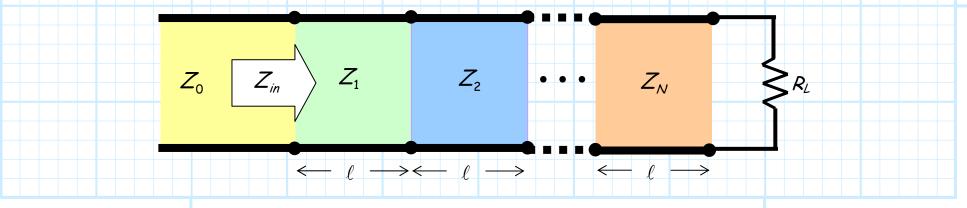
$$\Gamma_{in}(\omega) = \Gamma_0 + \Gamma_1 e^{-j2\omega T} + \Gamma_2 e^{-j4\omega T} + \dots + \Gamma_N e^{-j2N\omega T}$$
$$= \sum_{n=1}^{N} \Gamma_n e^{-j2n\omega T}$$

where:

$$\Gamma \doteq \frac{\ell}{\nu_p} = \text{propagation time through 1 section}$$

Note that for a multi-section transformer, we have N degrees of design freedom, corresponding to the N characteristic impedance values  $Z_n$ .

*n*=0



### **Behold the Binomial Function!**

**Q**: What should the values of  $\Gamma_n$  (i.e.,  $Z_n$ ) be?

A: We need to define Nindependent design equations, which we can then use to solve for the Nvalues of characteristic impedance  $Z_n$ .

First, we start with a single **design frequency**  $\omega_0$ , where we wish to achieve a **perfect** match:

$$\Gamma_{in}\left(\omega=\omega_{0}\right)=0$$

That's just one design equation: we need N - 1 more!

These addition equations can be selected using **many** criteria—one such criterion is to make the function  $\Gamma_{in}(\omega)$  **maximally flat** at the point  $\omega = \omega_0$ .

To accomplish this, we first consider the **Binomial Function**:

$$\Gamma(\theta) = \mathcal{A} \left( \mathbf{1} + \boldsymbol{e}^{-j2\theta} \right)^{\prime}$$

# What's so special about the Binomial Function!

The Binomial Function has the desirable properties that:

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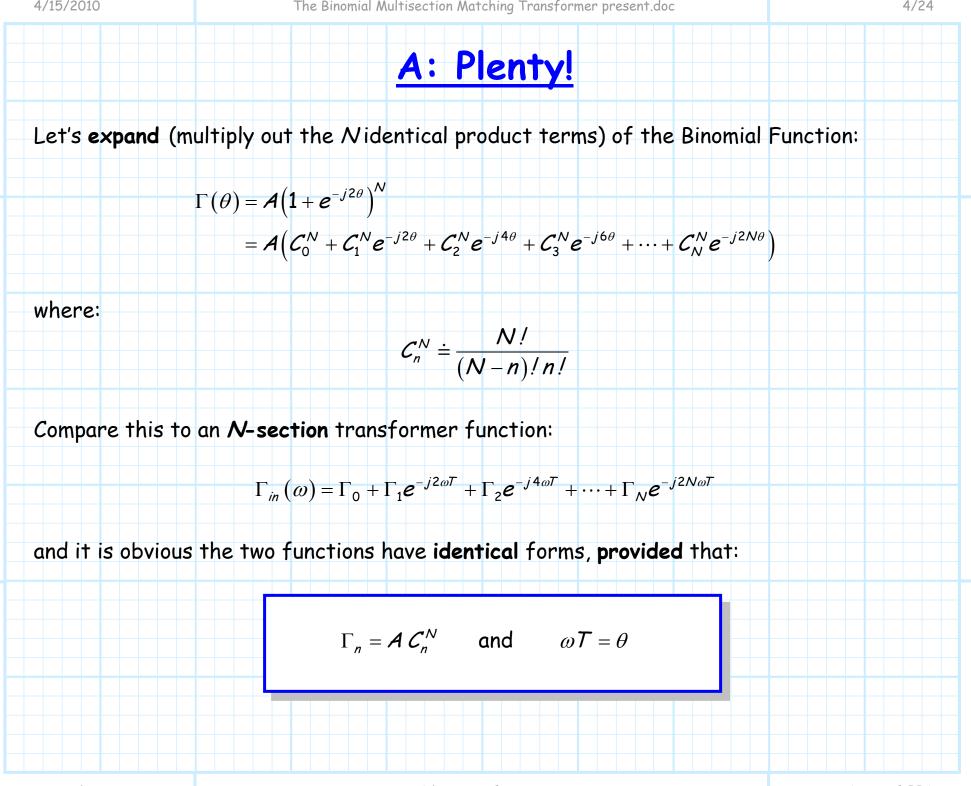
$$ig( heta = \pi/2 ig) = A ig( 1 + e^{-j\pi} ig)^N$$
  
=  $A ig( 1 - 1 ig)^N$   
=  $0$ 

and that:

$$\frac{d^n \Gamma(\theta)}{d\theta^n} \bigg|_{\theta=\frac{\pi}{2}} = 0 \text{ for } n = 1, 2, 3, \cdots, N-1$$

In other words, this Binomial Function is **maximally flat** at the point  $\theta = \pi/2$ , where it has a value of  $\Gamma(\theta = \pi/2) = 0$ .

Q: So? What does this have to do with our multi-section matching network?



### See, the Binomial Function is very useful!

Moreover, we find that this function is very **desirable** from the standpoint of a matching network. Recall that  $\Gamma(\theta) = 0$  at  $\theta = \pi/2$ --a **perfect** match!

Additionally, the function is maximally flat at  $\theta = \pi/2$ , therefore  $\Gamma(\theta) \approx 0$  over a wide range around  $\theta = \pi/2$ --a wide bandwidth!

**Q**: But how does  $\theta = \pi/2$  relate to frequency  $\omega$ ?

A: Remember that  $\omega T = \theta$ , so the value  $\theta = \pi/2$  corresponds to the frequency:

$$\omega_0 = \frac{1}{T} \frac{\pi}{2} = \frac{\nu_p}{\ell} \frac{\pi}{2}$$

This frequency  $(\omega_0)$  is therefore our **design** frequency—the frequency where we have a **perfect** match.

### What about the length of each section?

Note that the section-length  $\ell$  has an interesting **relationship** with this frequency:

$$P = \frac{\nu_p}{\omega_0} \frac{\pi}{2} = \frac{1}{\beta_0} \frac{\pi}{2} = \frac{\lambda_0}{2\pi} \frac{\pi}{2} = \frac{\lambda_0}{4}$$

In other words, a **Binomial** Multi-section matching network will have a **perfect** match at the frequency where the section lengths  $\ell$  are a **quarter wavelength**!

Thus, we have our first design rule:

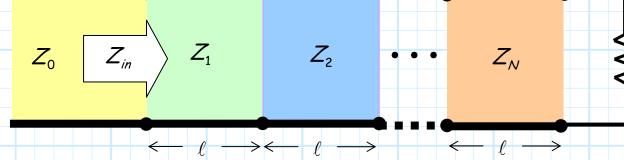
Set section lengths  $\ell$  so that they are a **quarter-wavelength**  $(\lambda_0/4)$  at the design frequency  $\omega_0$ .

### And that pesky constant A?

**Q**: I see! And then we select all the values  $Z_n$  such that  $\Gamma_n = A C_n^N$ . But wait! What is the value of **A** ??

A: We can determine this value by evaluating a boundary condition!

Specifically, we can **easily** determine the value of  $\Gamma(\omega)$  at  $\omega = 0$ .



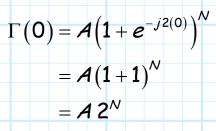
Note as  $\omega$  approaches **zero**, the electrical length  $\beta \ell$  of each section will **likewise** approach zero. Thus, the input impedance  $Z_{in}$  will simply be equal to  $R_L$  as  $\omega \to 0$ .

As a result, the input reflection coefficient  $\Gamma(\omega = 0)$  must be:

$$\Gamma(\omega=0) = \frac{Z_{in}(\omega=0) - Z_{0}}{Z_{in}(\omega=0) + Z_{0}} = \frac{R_{L} - Z_{0}}{R_{L} + Z_{0}}$$

# Aren't boundary conditions great ?

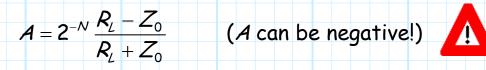
However, we likewise know that:



Equating the two expressions:

$$\Gamma(0) = A 2^{N} = \frac{R_{L} - Z_{0}}{R_{L} + Z_{0}}$$

And therefore:



We now have a form to calculate the required marginal reflection coefficients  $\Gamma_n$ :

$$\Gamma_n = AC_n^N = \frac{AN!}{(N-n)!n!}$$

### How do I determine characteristic impedance?

Of course, we **also** know that these marginal reflection coefficients are physically related to the characteristic impedances of each section as:

$$\Gamma_n = \frac{Z_{n+1} - Z_n}{Z_{n+1} + Z_n}$$

Equating the two and solving, we find that that the section characteristic impedances must satisfy:

$$Z_{n+1} = Z_n \frac{1+\Gamma_n}{1-\Gamma_n} = Z_n \frac{1+AC_n^N}{1-AC_n^N}$$

Note this is an **iterative** result—we determine  $Z_1$  from  $Z_0$ ,  $Z_2$  from  $Z_1$ , and so forth.

**Q:** This result **appears** to be our second design equation. Is there some reason why you didn't draw a big blue box around it?

A: Alas, there is a **big problem** with this result.

### The BIG problem with this result!

Note that there are N+1 coefficients  $\Gamma_n$  (i.e.,  $n \in \{0, 1, \dots, N\}$ ) in the Binomial series, yet there are only N design degrees of freedom (i.e., there are only N transmission line sections!).

Thus, our design is a bit over constrained, a result that manifests itself the finally marginal reflection coefficient  $\Gamma_N$ .

Note from the iterative solution above, the **last** transmission line impedance  $Z_N$  is selected to satisfy the **mathematical** requirement of the **penultimate** reflection coefficient  $\Gamma_{N-1}$ :

$$\Gamma_{N-1} = \frac{Z_{N} - Z_{N-1}}{Z_{N} + Z_{N-1}} = A C_{N-1}^{N}$$

Thus the last impedance must be:

$$Z_{N} = Z_{N-1} rac{1 + A C_{N-1}^{N}}{1 - A C_{N-1}^{N}}$$

#### Our degrees of freedom have run out!

But there is **one more** mathematical requirement! The last marginal reflection coefficient **must** likewise satisfy:

$$\Gamma_{N} = A C_{N}^{N} = 2^{-N} \frac{R_{L} - Z_{0}}{R_{L} + Z_{0}}$$

where we have used the fact that  $C_N^N = 1$ .

But, we just selected  $Z_N$  to satisfy the requirement for  $\Gamma_{N-1}$ ,—we have no physical design parameter to satisfy this last mathematical requirement!

As a result, we find to our great consternation that the last requirement is not satisfied:

$$\Gamma_{N} = \frac{R_{L} - Z_{N}}{R_{L} + Z_{N}} \neq A C_{N}^{N} \quad \text{IIIII}$$

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**Q:** Yikes! Does this mean that the resulting matching network will **not** have the desired Binomial frequency response?

A: That's exactly what it means!

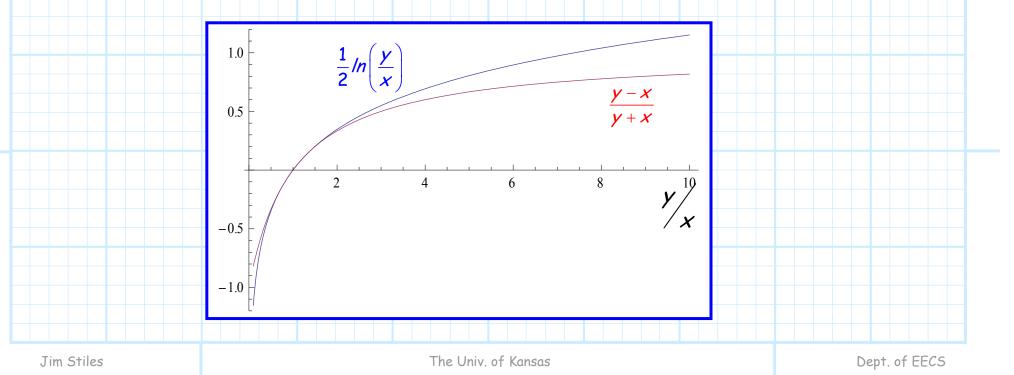
# <u>&\*#@\*!&!!!!!</u>

**Q:** You big #%@#\$%&!!!! **Why** did you **waste** all my time by discussing an overconstrained design problem that can't be built?

A: Relax; there is a solution to our dilemma—albeit an approximate one. You undoubtedly have previously used the approximation:

$$\frac{y-x}{y+x}\approx\frac{1}{2}\ln\left(\frac{y}{x}\right)$$

An approximation that is especially **accurate** when |y - x| is small (i.e., when  $\frac{y}{x} \simeq 1$ ).



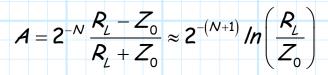
# Use this approximation for value A!

Now, we know that the values of  $Z_{n+1}$  and  $Z_n$  in a multi-section matching network are typically **very close**, such that  $|Z_{n+1} - Z_n|$  is small.

Thus, we use the approximation:

$$\Gamma_{n} = \frac{Z_{n+1} - Z_{n}}{Z_{n+1} + Z_{n}} \approx \frac{1}{2} \ln \left( \frac{Z_{n+1}}{Z_{n}} \right)$$

Likewise, we can **also** apply this approximation (although not as accurately) to the value of



A:

# Let's try this again—with approximations!

So, let's **start over**, only this time we'll use these **approximations**. First, determine A:

$$A \approx 2^{-(N+1)} ln \left(\frac{R_L}{Z_0}\right)$$

(A can be negative!)



Now use this result to calculate the **mathematically required** marginal reflection coefficients  $\Gamma_n$ :

$$\Gamma_n = \mathcal{A} \mathcal{C}_n^{\mathcal{N}} = \frac{\mathcal{A} \mathcal{N}!}{(\mathcal{N} - n)!n!}$$

# Here's (finally) our second design rule!

Of course, we **also** know that these marginal reflection coefficients are **physically** related to the **characteristic impedances** of each section as:

$$\Gamma_n \approx \frac{1}{2} \ln \left( \frac{Z_{n+1}}{Z_n} \right)$$

Equating the two and solving, we find that that the section characteristic impedances **must** satisfy:

$$Z_{n+1} = Z_n exp[2\Gamma_n]$$

Now this is our second design rule. Note it is an iterative rule—we determine  $Z_1$  from  $Z_0$ ,  $Z_2$  from  $Z_1$ , and so forth.

#### I don't understand what just happened

Q: Huh? How is this any better? How does applying **approximate** math lead to a **better** design result??

A: Applying these approximations help **resolve** our over-constrained problem. Recall that the over-constraint resulted in:

$$\Gamma_{N} = \frac{R_{L} - Z_{N}}{R_{L} + Z_{N}} \neq A C_{N}^{N}$$

But, as it turns out, these approximations leads to the happy situation where:

$$\Gamma_N \approx \frac{1}{2} ln \left( \frac{R_L}{Z_N} \right) = A C_N^N \quad \leftarrow \quad \text{A Sanity check!!}$$

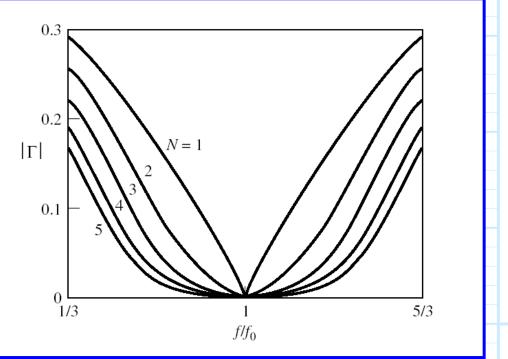
provided that the value A is likewise the approximation given above.

# I still don't understand what just happened

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Effectively, these approximations couple the results, such that each value of characteristic impedance  $Z_n$  approximately satisfies both  $\Gamma_n$  and  $\Gamma_{n+1}$ . Summarizing:

- \* If you use the "exact" design equations to determine the characteristic impedances  $Z_n$ , the last value  $\Gamma_N$  will exhibit a significant numeric error, and your design will not appear to be maximally flat.
- If you instead use the "approximate" design equations to determine the characteristic impedances Z<sub>n</sub>, all values Γ<sub>n</sub> will exhibit a slight error, but the resulting design will appear to be maximally flat, Binomial reflection coefficient function Γ(ω)!



Note that as we **increase** the number of **sections**, the matching **bandwidth** increases.

# Bandwidth: How do we define it?

**band**·width (band'width', -witth') - noun

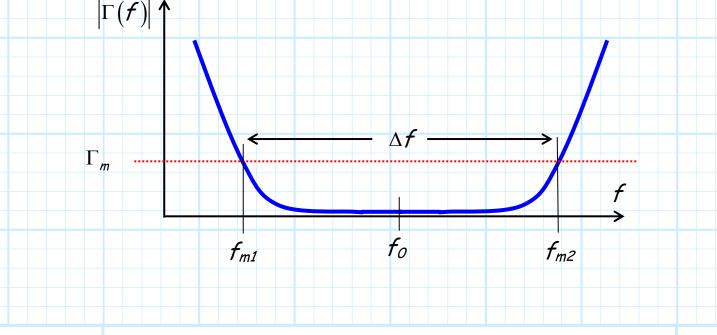
1. the range of frequencies within a ....



Can we determine the value of this bandwidth?

A: Sure! But we first must define what we mean by bandwidth.

As we move from the design (perfect match) frequency  $f_0$  the value  $|\Gamma(f)|$  will increase. At some frequency ( $f_m$ , say) the magnitude of the reflection coefficient will increase to some unacceptably high value ( $\Gamma_m$ , say). At that point, we no longer consider the device to be matched.



# Bandwidth: How do we calculate it?

Note there are two values of frequency  $f_m$  —one value less than design frequency  $f_0$ , and one value greater than design frequency  $f_0$ .

These two values define the **bandwidth**  $\Delta f$  of the matching network:

$$\Delta f = f_{m2} - f_{m1} = 2(f_0 - f_{m1}) = 2(f_{m2} - f_0)$$

**Q**: So what is the **numerical** value of  $\Gamma_m$ ?

A: I don't know—it's up to you to decide!

Every engineer must determine what they consider to be an acceptable match (i.e., decide

This decision depends on the **application** involved, and the **specifications** of the overall microwave system being designed.

However, we **typically** set  $\Gamma_m$  to be 0.2 or less.



 $\Gamma_m$ ).

 $f_m$ ?

# We get to perform some Algebra!!

**Q:** OK, after we have selected  $\Gamma_m$ , can we determine the **two** frequencies

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A: Sure! We just have to do a little algebra.

We start by **rewriting** the Binomial function:

$$\Gamma(\theta) = \mathcal{A} \left( 1 + e^{-j2\theta} \right)^{N}$$
$$= \mathcal{A} e^{-jN\theta} \left( e^{+j\theta} + e^{-j\theta} \right)^{N}$$
$$= \mathcal{A} e^{-jN\theta} \left( e^{+j\theta} + e^{-j\theta} \right)^{N}$$
$$= \mathcal{A} e^{-jN\theta} \left( 2\cos\theta \right)^{N}$$

Now, we take the **magnitude** of this function:

$$\Gamma(\theta) = 2^{N} |\mathbf{A}| |e^{-jN\theta}| |\cos\theta|^{N}$$

 $= 2^{N} |A| |\cos \theta|^{N}$ 

#### It gets better—even more algebra!!

Now, we **define** the values  $\theta$  where  $|\Gamma(\theta)| = \Gamma_m$  as  $\theta_m$ . I.E., :

$$\Gamma_{m} = \left| \Gamma \left( \theta = \theta_{m} \right) \right|$$
$$= 2^{N} \left| \mathcal{A} \right| \left| \cos \theta_{m} \right|^{N}$$

We can now solve for  $\theta_m$  (in **radians**!) in terms of  $\Gamma_m$ :

Note that there are **two solutions** to the above equation (one less that  $\pi/2$  and one greater than  $\pi/2$ )!

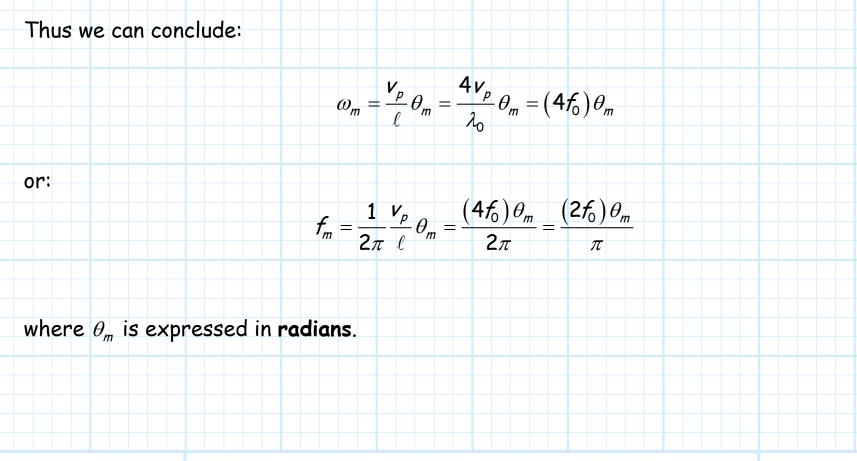
Now, we can **convert** the values of  $\theta_m$  into specific **frequencies**.

# Converting $\theta_m$ to $f_m$

Recall that  $\omega T = \theta$ , therefore:

$$\omega_m = \frac{1}{T} \theta_m = \frac{V_p}{\ell} \theta_m$$

But recall also that  $\ell = \lambda_0/4$ , where  $\lambda_0$  is the wavelength at the **design frequency**  $f_0$  (not  $f_m!$ ), and where  $\lambda_0 = v_p/f_0$ .



# And thus the bandwidth is...

Therefore:

$$f_{m1} = \frac{2f_0}{\pi} \cos^{-1} \left[ +\frac{1}{2} \left( \frac{\Gamma_m}{|\mathcal{A}|} \right)^{1/N} \right] \qquad f_{m2} = \frac{2f_0}{\pi} \cos^{-1} \left[ -\frac{1}{2} \left( \frac{\Gamma_m}{|\mathcal{A}|} \right)^{1/N} \right]$$

Thus, the bandwidth of the binomial matching network can be determined as:

$$\Delta f = 2(f_0 - f_{m1})$$
$$= 2f_0 - \frac{4f_0}{\pi} \cos^{-1} \left[ + \frac{1}{2} \left( \frac{\Gamma_m}{|\mathcal{A}|} \right)^{1/N} \right]$$

Note that this equation can be used to determine the **bandwidth** of a binomial matching network, given  $\Gamma_m$  and number of sections N.

However, it can likewise be used to determine the **number of sections** Nrequired to meet a specific bandwidth requirement!

#### In summary, our design steps

Finally, we can list the **design steps** for a binomial matching network:

- **1.** Determine the value N required to meet the bandwidth ( $\Delta f$  and  $\Gamma_m$ ) requirements.
- **2.** Determine the **approximate** value A from  $Z_0, R_1$  and N.
- **3**. Determine the marginal reflection coefficients  $\Gamma_n = AC_n^N$  required by the binomial function.
- 4. Determine the characteristic impedance of each section using the **iterative approximation**:

$$Z_{n+1} = Z_n \exp\left[2\Gamma_n\right]$$

5. Perform the sanity check:

$$\Gamma_{N} \approx \frac{1}{2} \ln \left( \frac{R_{L}}{Z_{N}} \right) = \mathcal{A} C_{N}^{N}$$

6. Determine section length  $\ell = \lambda_0/4$  for design frequency  $f_0$ .