The Binomial Multi-Section Transformer

Recall that a multi-section matching network can be described using the theory of small reflections as:

$$\Gamma_{in}(\omega) = \Gamma_0 + \Gamma_1 e^{-j2\omega T} + \Gamma_2 e^{-j4\omega T} + \dots + \Gamma_N e^{-j2N\omega T}$$
$$= \sum_{n=0}^{N} \Gamma_n e^{-j2n\omega T}$$

where:

$$T \doteq \frac{\ell}{\nu_p}$$
 = propagation time through 1 section

Note that for a multi-section transformer, we have N degrees of design freedom, corresponding to the N characteristic impedance values Z_n .

Q: What should the values of Γ_n (i.e., Z_n) be?

A: We need to define Nindependent design equations, which we can then use to solve for the Nvalues of characteristic impedance Z_n .

First, we start with a single design frequency ω_0 , where we wish to achieve a perfect match:

$$\Gamma_{in}(\omega=\omega_0)=0$$

That's just one design equation: we need N-1 more!

These addition equations can be selected using many criteria—one such criterion is to make the function $\Gamma_{in}(\omega)$ maximally flat at the point $\omega = \omega_0$.

To accomplish this, we first consider the Binomial Function:

$$\Gamma(\theta) = A(1 + e^{-j2\theta})^N$$

This function has the desirable properties that:

$$\Gamma(\theta = \pi/2) = A(1 + e^{-j\pi})^{N}$$
$$= A(1-1)^{N}$$
$$= 0$$

and that:

$$\frac{d^n \Gamma(\theta)}{d\theta^n}\bigg|_{\theta=\frac{\pi}{2}} = 0 \text{ for } n = 1, 2, 3, \dots, N-1$$

In other words, this Binomial Function is **maximally flat** at the point $\theta = \pi/2$, where it has a value of $\Gamma(\theta = \pi/2) = 0$.

Q: So? What does this have to do with our multi-section matching network?

A: Let's expand (multiply out the Nidentical product terms) of the Binomial Function:

$$\Gamma(\theta) = A \left(1 + e^{-j2\theta} \right)^{N}$$

$$= A \left(C_{0}^{N} + C_{1}^{N} e^{-j2\theta} + C_{2}^{N} e^{-j4\theta} + C_{3}^{N} e^{-j6\theta} + \dots + C_{N}^{N} e^{-j2N\theta} \right)$$

where:

$$C_n^N \doteq \frac{N!}{(N-n)! \, n!}$$

Compare this to an N-section transformer function:

$$\Gamma_{in}(\omega) = \Gamma_0 + \Gamma_1 e^{-j2\omega T} + \Gamma_2 e^{-j4\omega T} + \dots + \Gamma_N e^{-j2N\omega T}$$

and it is obvious the two functions have identical forms, provided that:

$$\Gamma_n = A C_n^N$$
 and $\omega T = \theta$

Moreover, we find that this function is very **desirable** from the standpoint of the a matching network. Recall that $\Gamma(\theta) = 0$ at $\theta = \pi/2$ --a **perfect** match!

Additionally, the function is **maximally flat** at $\theta=\pi/2$, therefore $\Gamma(\theta)\approx 0$ over a wide range around $\theta=\pi/2$ --a wide bandwidth!

Q: But how does $\theta = \pi/2$ relate to frequency ω ?

A: Remember that $\omega T = \theta$, so the value $\theta = \pi/2$ corresponds to the frequency:

$$\omega_0 = \frac{1}{T} \frac{\pi}{2} = \frac{V_p}{\ell} \frac{\pi}{2}$$

This frequency (ω_0) is therefore our **design** frequency—the frequency where we have a **perfect** match.

Note that the length ℓ has an interesting **relationship** with this frequency:

$$\ell = \frac{v_p}{\omega_0} \frac{\pi}{2} = \frac{1}{\beta_0} \frac{\pi}{2} = \frac{\lambda_0}{2\pi} \frac{\pi}{2} = \frac{\lambda_0}{4}$$

In other words, a Binomial Multi-section matching network will have a perfect match at the frequency where the section lengths ℓ are a quarter wavelength!

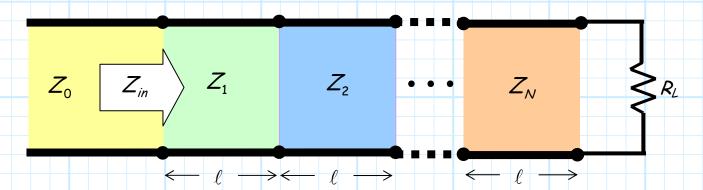
Thus, we have our first design rule:

Set section lengths ℓ so that they are a quarter-wavelength $(\lambda_0/4)$ at the design frequency ω_0 .

Q: I see! And then we select all the values Z_n such that $\Gamma_n = A C_n^N$. But wait! What is the value of A??

A: We can determine this value by evaluating a boundary condition!

Specifically, we can **easily** determine the value of $\Gamma(\omega)$ at $\omega = 0$.



Note as ω approaches **zero**, the electrical length $\beta \ell$ of each section will **likewise** approach zero. Thus, the input impedance Z_{in} will simply be equal to R_{ℓ} as $\omega \to 0$.

As a result, the input reflection coefficient $\Gamma(\omega=0)$ must be:

$$\Gamma(\omega = 0) = \frac{Z_{in}(\omega = 0) - Z_0}{Z_{in}(\omega = 0) + Z_0}$$
$$= \frac{R_L - Z_0}{R_L + Z_0}$$

However, we likewise know that:

$$\Gamma(0) = A(1 + e^{-j2(0)})^{N}$$

$$= A(1+1)^{N}$$

$$= A 2^{N}$$

Equating the two expressions:

$$\Gamma(0) = A 2^{N} = \frac{R_{L} - Z_{0}}{R_{L} + Z_{0}}$$

And therefore:

$$A = 2^{-N} \frac{R_L - Z_0}{R_L + Z_0}$$
 (A can be negative!)



We now have a form to calculate the required marginal reflection coefficients Γ_n :

$$\Gamma_n = AC_n^N = \frac{AN!}{(N-n)!n!}$$

Of course, we also know that these marginal reflection coefficients are physically related to the characteristic impedances of each section as:

$$\Gamma_n = \frac{Z_{n+1} - Z_n}{Z_{n+1} + Z_n}$$

Equating the two and solving, we find that that the section characteristic impedances must satisfy:

$$Z_{n+1} = Z_n \frac{1+\Gamma_n}{1-\Gamma_n} = Z_n \frac{1+AC_n^N}{1-AC_n^N}$$

Note this is an **iterative** result—we determine Z_1 from Z_0 , Z_2 from Z_1 , and so forth.

Q: This result **appears** to be our second design equation. Is there some reason why you didn't draw a big blue box around it?

A: Alas, there is a big problem with this result.

Note that there are N+1 coefficients Γ_n (i.e., $n \in \{0,1,\dots,N\}$) in the Binomial series, yet there are only N design degrees of freedom (i.e., there are only N transmission line sections!).

Thus, our design is a bit over constrained, a result that manifests itself the finally marginal reflection coefficient $\Gamma_{\mathcal{N}}$.

Note from the iterative solution above, the **last** transmission line impedance Z_N is selected to satisfy the **mathematical** requirement of the **penultimate** reflection coefficient Γ_{N-1} :

$$\Gamma_{N-1} = \frac{Z_N - Z_{N-1}}{Z_N + Z_{N-1}} = A C_{N-1}^N$$

Thus the last impedance must be:

$$Z_N = Z_{N-1} \frac{1 + A C_{N-1}^N}{1 - A C_{N-1}^N}$$

But there is **one more** mathematical requirement! The last marginal reflection coefficient **must** likewise satisfy:

$$\Gamma_{N} = A C_{N}^{N} = 2^{-N} \frac{R_{L} - Z_{0}}{R_{L} + Z_{0}}$$

where we have used the fact that $C_N^N = 1$.

But, we just selected Z_N to satisfy the requirement for Γ_{N-1} ,—we have no **physical** design parameter to satisfy this last **mathematical** requirement!

As a result, we find to our great consternation that the last requirement is not satisfied:

$$\Gamma_{N} = \frac{R_{L} - Z_{N}}{R_{L} + Z_{N}} \neq A C_{N}^{N}$$
 ||||||

Q: Yikes! Does this mean that the resulting matching network will not have the desired Binomial frequency response?

A: That's exactly what it means!

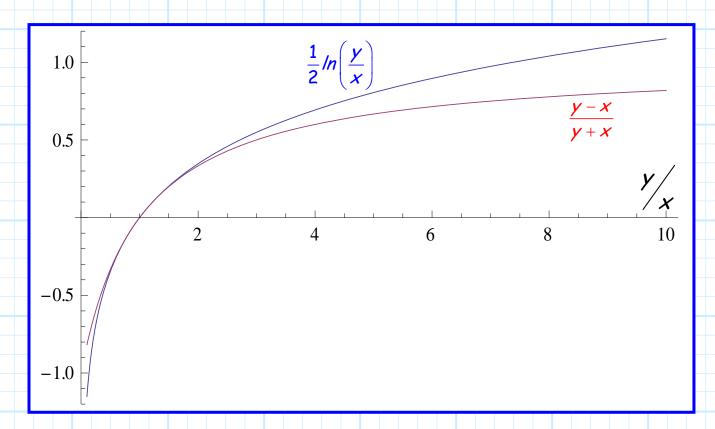
Q: You big #%@#\$%&!!!! Why did you waste all my time by discussing an over-constrained design problem that can't be built?

A: Relax; there is a solution to our dilemma—albeit an approximate one.

You undoubtedly have previously used the approximation:

$$\frac{y-x}{y+x} \approx \frac{1}{2} \ln \left(\frac{y}{x} \right)$$

An approximation that is especially accurate when |y-x| is small (i.e., when $\frac{y}{x} \simeq 1$).



Now, we know that the values of Z_{n+1} and Z_n in a multi-section matching network are typically **very close**, such that $\left|Z_{n+1}-Z_n\right|$ is small. Thus, we use the approximation:

$$\Gamma_n = \frac{Z_{n+1} - Z_n}{Z_{n+1} + Z_n} \approx \frac{1}{2} \ln \left(\frac{Z_{n+1}}{Z_n} \right)$$

Likewise, we can also apply this approximation (although not as accurately) to the value of A:

$$A = 2^{-N} \frac{R_L - Z_0}{R_L + Z_0} \approx 2^{-(N+1)} \ln \left(\frac{R_L}{Z_0}\right)$$

So, let's start over, only this time we'll use these approximations. First, determine A:

$$A \approx 2^{-(N+1)} \ln \left(\frac{R_L}{Z_0} \right)$$
 (A can be negative!)



Now use this result to calculate the mathematically required marginal reflection coefficients Γ_n :

$$\Gamma_n = AC_n^N = \frac{AN!}{(N-n)!n!}$$

Of course, we also know that these marginal reflection coefficients are physically related to the characteristic impedances of each section as:

$$\Gamma_n \approx \frac{1}{2} ln \left(\frac{Z_{n+1}}{Z_n} \right)$$

Equating the two and solving, we find that that the section characteristic impedances **must** satisfy:

$$Z_{n+1} = Z_n \exp[2\Gamma_n]$$

Now this is our second design rule. Note it is an iterative rule—we determine Z_1 from Z_0 , Z_2 from Z_1 , and so forth.

Q: Huh? How is this any better? How does applying approximate math lead to a better design result??

A: Applying these approximations help resolve our overconstrained problem. Recall that the over-constraint resulted in:

$$\Gamma_{N} = \frac{R_{L} - Z_{N}}{R_{L} + Z_{N}} \neq A C_{N}^{N}$$

But, as it turns out, these approximations leads to the happy situation where:

$$\Gamma_N \approx \frac{1}{2} ln \left(\frac{R_L}{Z_N} \right) = A C_N^N \quad \leftarrow \text{Sanity check!!}$$

provided that the value A is likewise the approximation given above.

Effectively, these approximations couple the results, such that each value of characteristic impedance Z_n approximately satisfies both Γ_n and Γ_{n+1} . Summarizing:

- * If you use the "exact" design equations to determine the characteristic impedances Z_n , the last value Γ_N will exhibit a significant numeric error, and your design will not appear to be maximally flat.
- * If you instead use the "approximate" design equations to determine the characteristic impedances Z_n , all values Γ_n will exhibit a slight error, but the resulting design will appear to be maximally flat, Binomial reflection coefficient function $\Gamma(\omega)!$

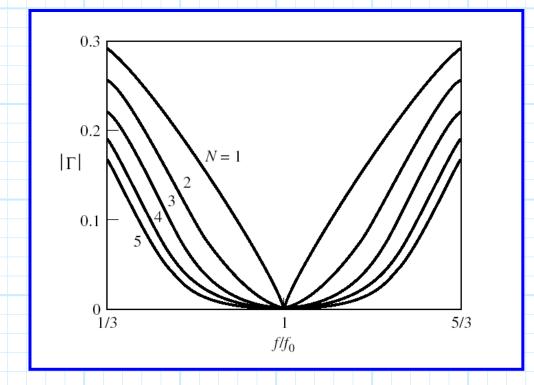


Figure 5.15 (p. 250)

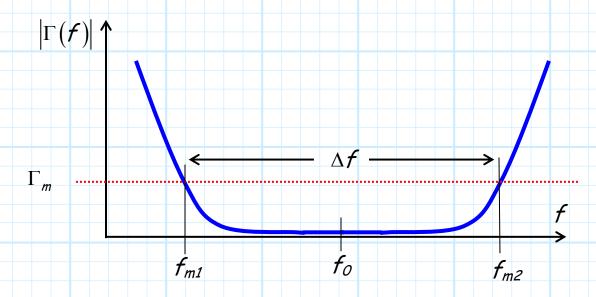
Reflection coefficient magnitude versus frequency for multisection binomial matching transformers of Example 5.6 Z_L = 50 Ω and Z_0 = 100 Ω .

Note that as we increase the number of sections, the matching bandwidth increases.

Q: Can we determine the value of this bandwidth?

A: Sure! But we first must define what we mean by bandwidth.

As we move from the design (perfect match) frequency f_0 the value $|\Gamma(f)|$ will **increase**. At some frequency (f_m, say) the magnitude of the reflection coefficient will increase to some **unacceptably** high value (Γ_m, say) . At that point, we **no longer** consider the device to be matched.



Note there are **two** values of frequency f_m —one value **less** than design frequency f_O , and one value **greater** than design frequency f_O . These two values define the **bandwidth** Δf of the matching network:

$$\Delta f = f_{m2} - f_{m1} = 2(f_0 - f_{m1}) = 2(f_{m2} - f_0)$$

Q: So what is the numerical value of Γ_m ?

A: I don't know—it's up to you to decide!

Every engineer must determine what **they** consider to be an acceptable match (i.e., decide Γ_m). This decision depends on the **application** involved, and the **specifications** of the overall microwave system being designed.

However, we **typically** set Γ_m to be 0.2 or less.

Q: OK, after we have selected Γ_m , can we determine the **two** frequencies f_m ?

A: Sure! We just have to do a little algebra.

We start by rewriting the Binomial function:

$$\Gamma(\theta) = A \left(1 + e^{-j2\theta} \right)^{N}$$

$$= A e^{-jN\theta} \left(e^{+j\theta} + e^{-j\theta} \right)^{N}$$

$$= A e^{-jN\theta} \left(e^{+j\theta} + e^{-j\theta} \right)^{N}$$

$$= A e^{-jN\theta} \left(2\cos\theta \right)^{N}$$

Now, we take the **magnitude** of this function:

$$|\Gamma(\theta)| = 2^{N} |A| |e^{-jN\theta}| |\cos\theta|^{N}$$
$$= 2^{N} |A| |\cos\theta|^{N}$$

Now, we define the values θ where $|\Gamma(\theta)| = \Gamma_m$ as θ_m . I.E., :

$$\Gamma_{m} = \left| \Gamma(\theta = \theta_{m}) \right|$$
$$= 2^{N} |A| |\cos \theta_{m}|^{N}$$

We can now solve for θ_m (in radians!) in terms of Γ_m :

$$\theta_{m1} = \cos^{-1} \left[\frac{1}{2} \left(\frac{\Gamma_m}{|A|} \right)^{\frac{1}{N}} \right] \qquad \theta_{m2} = \cos^{-1} \left[-\frac{1}{2} \left(\frac{\Gamma_m}{|A|} \right)^{\frac{1}{N}} \right]$$

Note that there are **two solutions** to the above equation (one less that $\pi/2$ and one greater than $\pi/2$)!

Now, we can convert the values of θ_m into specific frequencies.

Recall that $\omega T = \theta$, therefore:

$$\omega_m = \frac{1}{T} \theta_m = \frac{V_p}{\ell} \theta_m$$

But recall also that $\ell=\lambda_0/4$, where λ_0 is the wavelength at the **design frequency** f_0 (not $f_m!$), and where $\lambda_0=v_p/f_0$.

Thus we can conclude:

$$\omega_m = \frac{v_p}{\ell} \theta_m = \frac{4v_p}{\lambda_0} \theta_m = (4f_0) \theta_m$$

or:

$$f_{m} = \frac{1}{2\pi} \frac{v_{p}}{\ell} \theta_{m} = \frac{(4f_{0}) \theta_{m}}{2\pi} = \frac{(2f_{0}) \theta_{m}}{\pi}$$

where θ_m is expressed in radians. Therefore:

$$f_{m1} = \frac{2f_0}{\pi} \cos^{-1} \left[+ \frac{1}{2} \left(\frac{\Gamma_m}{|\mathcal{A}|} \right)^{1/N} \right] \qquad f_{m2} = \frac{2f_0}{\pi} \cos^{-1} \left[- \frac{1}{2} \left(\frac{\Gamma_m}{|\mathcal{A}|} \right)^{1/N} \right]$$

Thus, the **bandwidth** of the binomial matching network can be determined as:

$$\Delta f = 2(f_0 - f_{m1})$$

$$= 2f_0 - \frac{4f_0}{\pi} \cos^{-1} \left[+ \frac{1}{2} \left(\frac{\Gamma_m}{|A|} \right)^{1/N} \right]$$

Note that this equation can be used to determine the **bandwidth** of a binomial matching network, given Γ_m and number of sections N.

However, it can likewise be used to determine the number of sections Nrequired to meet a specific bandwidth requirement!

Finally, we can list the **design steps** for a binomial matching network:

- 1. Determine the value N required to meet the bandwidth (Δf and Γ_m) requirements.
- 2. Determine the approximate value A from Z_0 , R_L and N.
- 3. Determine the marginal reflection coefficients $\Gamma_n = A C_n^N$ required by the binomial function.
- 4. Determine the characteristic impedance of each section using the iterative approximation:

$$Z_{n+1} = Z_n \exp[2\Gamma_n]$$

5. Perform the sanity check:

$$\Gamma_N \approx \frac{1}{2} \ln \left(\frac{R_L}{Z_N} \right) = A C_N^N$$

6. Determine section length $\ell = \lambda_0/4$ for design frequency f_0 .