5.6 – Binomial Multi-section Matching Transformer

Reading Assignment: *pp. 246-250*

One way to **maximize bandwidth** is to construct a multisection matching network with a function $\Gamma(f)$ that is **maximally flat**.

Q: *Maximally flat? What kind of function is **maximally flat?***

This function maximizes bandwidth by providing a solution that is **maximally flat**.

A: HO: MAXIMALLY FLAT FUNCTIONS

1. We can build a multisection matching network such that the function $\Gamma(f)$ is a **binomial function**.

2. The binomial function is **maximally flat**.

Q: *Meaning?*

A: Meaning the function $\Gamma(f)$ is maximally flat→ a **wideband** solution!

HO: THE BINOMIAL MULTI-SECTION MATCHING TRANSFORMER

Maximally Flat Functions

Consider some function $f(x)$. Say that we know the value of the function **at** $x=1$ is 5:

$$f(x=1)=5$$

This of course says **something** about the function $f(x)$, but it **doesn't** tell us much!

We can additionally determine the **first derivative** of this function, and likewise evaluate this derivative **at** $x=1$. Say that this value turns out to be **zero**:

$$\left.\frac{d\,f(x)}{d\,x}\right|_{x=1}=0$$

Note that this does not mean that the derivative of $f(x)$ is equal to zero, it merely means that the derivative of $f(x)$ is zero **at the value** $x=1$. Presumably, $d\,f(x)/dx$ is **non-zero** at **other** values of x.

Taking derivates: way too fun to stop!

So, we now have **two** pieces of information about the function $f(x)$. We can add to this list by continuing to take higher-order derivatives and evaluating them at the single point $x=1$.

Let's say that the values of **all** the derivatives (at $x=1$) turn out to have a zero value:

$$\left.\frac{d f^n(x)}{d x^n}\right|_{x=1} = 0 \quad \text{for } n = 1,2,3,\cdots,\infty$$

We say that this function is **completely flat** at the point x=1.

Because **all** the derivatives are zero at $x=1$, it means that the function cannot change in value from that at $x=1$.

In other words, if the function has a value of 5 at $x=1$, (i.e., $f(x=1)=5$), then the function **must** have a value of 5 at **all** x !

The function $f(x)$ thus must be the **constant** function:

$$f(x) = 5$$

A more realistic function

Now let's consider the following **problem**—say some function $f(x)$ has the following form:

$$f(x) = a x^3 + b x^2 + c x$$

We wish to **determine** the values a, b, and c so that:

$$f(x=1) = 5$$

and that the value of the function $f(x)$ is as **close** to a value of 5 as possible in the region where x =1.

In other words, we want the function to have the value of 5 at x =1, and to **change** from that value as **slowly** as possible as we "move" from x =1.

Completely flat in not possible!

Q: *Don't we simply want the* ***completely*** *flat function* $f(x)=5$*?*

A: That would be the **ideal** function for this case, but notice that solution is **not** an option. Note there are **no** values of *a*, *b*, and *c* that will make:

$$a x^3 + b x^2 + c x = 5$$

for **all** values x.

Q: *So* ***what*** *do we do?*

A: **Instead** of the completely flat solution, we can find the **maximally flat** solution!

The **maximally flat** solution comes from determining the values *a*, *b*, and *c* so that as many derivatives **as possible** are **zero** at the point x=1.

How many derivatives can be zero?

For example, we wish to make the **first derivate** equal to zero at $x=1$:

$$
\begin{aligned}
0 &= \left. \frac{d f(x)}{d x} \right|_{x=1} \\
&= \left. \left(3ax^2 + 2bx + c\right) \right|_{x=1} \\
&= 3a + 2b + c
\end{aligned}
$$

Likewise, we wish to make the **second derivative** equal to zero at $x=1$:

$$
\begin{aligned}
0 &= \left. \frac{d^2 f(x)}{d x^2} \right|_{x=1} \\
&= \left. \left(6ax + 2b\right) \right|_{x=1} \\
&= 6a + 2b
\end{aligned}
$$

Here we must **stop** taking derivatives, as our solution only has **three degrees of design freedom** (i.e., 3 unknowns a, b, c).

We're out of degrees of design freedom

Q: *But we only have taken **two** derivatives, can't we take **one more**?*

A: **No!** We already have a **third** "design" equation: the value of the function **must** be 5 at $x = 1$:

$$\begin{aligned} 5 &= f(x=1) \\ &= a(1)^3 + b(1)^2 + c(1) \\ &= a + b + c \end{aligned}$$

So, we have used the **maximally flat** criterion at $x = 1$ to generate **three** equations and **three** unknowns:

$$5 = a + b + c$$

$$0 = 3a + 2b + c$$

$$0 = 6a + 2b$$

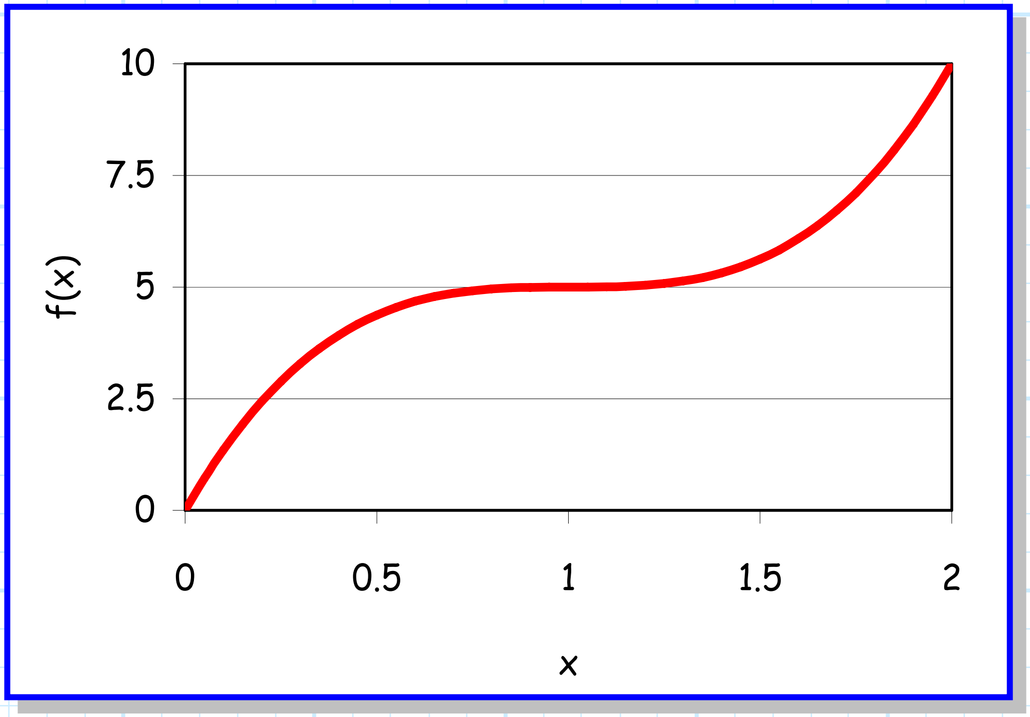
Solving, we find:

$$\begin{aligned} a &= 5 \\ b &= -15 \\ c &= 15 \end{aligned}$$

Look! The function is maximally flat at x=1!

Therefore, the **maximally flat** function (at $x=1$) is:

$$f(x) = 5x^3 - 15x^2 + 15x$$



The Binomial Multi-Section Transformer

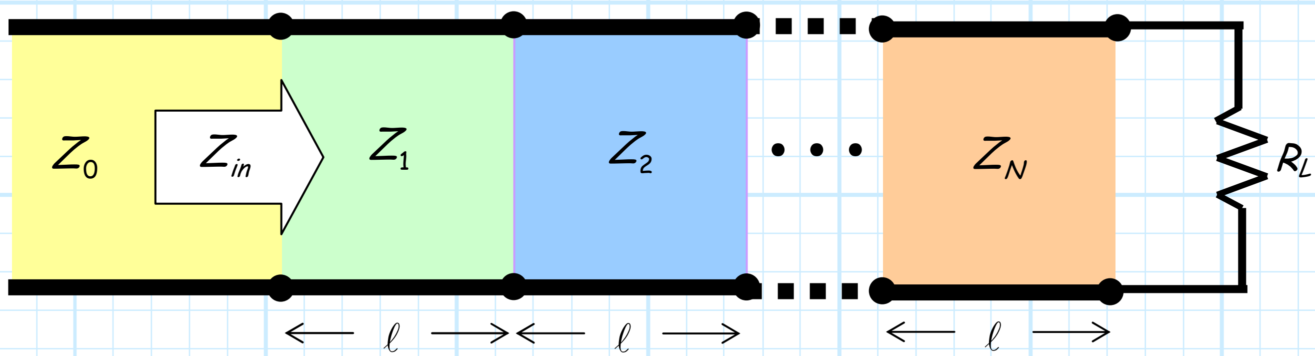
Recall that a **multi-section matching network** can be described using the theory of small reflections as:

$$\begin{aligned}\Gamma_{in}(\omega) &= \Gamma_0 + \Gamma_1 e^{-j2\omega T} + \Gamma_2 e^{-j4\omega T} + \cdots + \Gamma_N e^{-j2N\omega T} \\ &= \sum_{n=0}^{N} \Gamma_n e^{-j2n\omega T}\end{aligned}$$

where:

$$T \doteq \frac{\ell}{v_p} = \text{propagation time through 1 section}$$

Note that for a multi-section transformer, we have N **degrees of design freedom**, corresponding to the N characteristic impedance values Z_n.



Behold the Binomial Function!

Q: *What should the values of Γ_n (i.e., Z_n) be?*

A: We need to define N independent **design equations**, which we can then use to solve for the N values of **characteristic impedance** Z_n.

First, we start with a single **design frequency** ω_0, where we wish to achieve a **perfect** match:

$$\Gamma_{in}\left(\omega = \omega_0\right) = 0$$

That's just **one** design equation: we need $\boldsymbol{N - 1}$ more!

These addition equations can be selected using **many** criteria—one such criterion is to make the function $\Gamma_{in}(\omega)$ **maximally flat** at the point $\omega = \omega_0$.

To accomplish this, we first consider the **Binomial Function**:

$$\Gamma(\theta) = A\left(1 + e^{-j2\theta}\right)^N$$

What's so special about the Binomial Function!

The Binomial Function has the desirable **properties** that:

$$\begin{aligned}\Gamma\left(\theta=\pi/2\right) &= A\left(1+e^{-j\pi}\right)^N \\ &= A(1-1)^N \\ &= 0\end{aligned}$$

and that:

$$\left.\frac{d^n\,\Gamma(\theta)}{d\theta^n}\right|_{\theta=\pi/2} = 0 \text{ for } n=1,2,3,\cdots,N-1$$

In other words, this Binomial Function is **maximally flat** at the point $\theta=\pi/2$, where it has a value of $\Gamma\left(\theta=\pi/2\right)=0$.

Q: *So? What does **this** have to do with our **multi-section matching network**?*

A: Plenty!

Let's **expand** (multiply out the *N* identical product terms) of the Binomial Function:

$$\begin{aligned}\Gamma(\theta) &= A\left(1+e^{-j2\theta}\right)^N \\ &= A\left(C_0^N + C_1^N e^{-j2\theta} + C_2^N e^{-j4\theta} + C_3^N e^{-j6\theta} + \cdots + C_N^N e^{-j2N\theta}\right)\end{aligned}$$

where:

$$C_n^N \doteq \frac{N!}{(N-n)!\,n!}$$

Compare this to an ***N*-section** transformer function:

$$\Gamma_{in}(\omega) = \Gamma_0 + \Gamma_1 e^{-j2\omega T} + \Gamma_2 e^{-j4\omega T} + \cdots + \Gamma_N e^{-j2N\omega T}$$

and it is obvious the two functions have **identical** forms, **provided** that:

$$\Gamma_n = A\,C_n^N \quad \text{and} \quad \omega T = \theta$$

See, the Binomial Function is very useful!

Moreover, we find that this function is very **desirable** from the standpoint of a matching network. Recall that $\Gamma(\theta) = 0$ at $\theta = \pi/2$--a **perfect** match!

Additionally, the function is **maximally flat** at $\theta = \pi/2$, therefore $\Gamma(\theta) \approx 0$ over a wide range around $\theta = \pi/2$--a **wide bandwidth**!

Q: *But how does $\theta = \pi/2$ relate to frequency ω?*

A: Remember that $\omega T = \theta$, so the value $\theta = \pi/2$ corresponds to the frequency:

$$\omega_0 = \frac{1}{T}\frac{\pi}{2} = \frac{v_p}{\ell}\frac{\pi}{2}$$

This frequency (ω_0) is therefore our **design** frequency—the frequency where we have a **perfect** match.

What about the length of each section?

Note that the section-length ℓ has an interesting **relationship** with this frequency:

$$\ell = \frac{v_p}{\omega_0}\frac{\pi}{2} = \frac{1}{\beta_0}\frac{\pi}{2} = \frac{\lambda_0}{2\pi}\frac{\pi}{2} = \frac{\lambda_0}{4}$$

In other words, a **Binomial** Multi-section matching network will have a **perfect** match at the frequency where the section lengths ℓ are a **quarter wavelength**!

Thus, we have our **first design rule**:

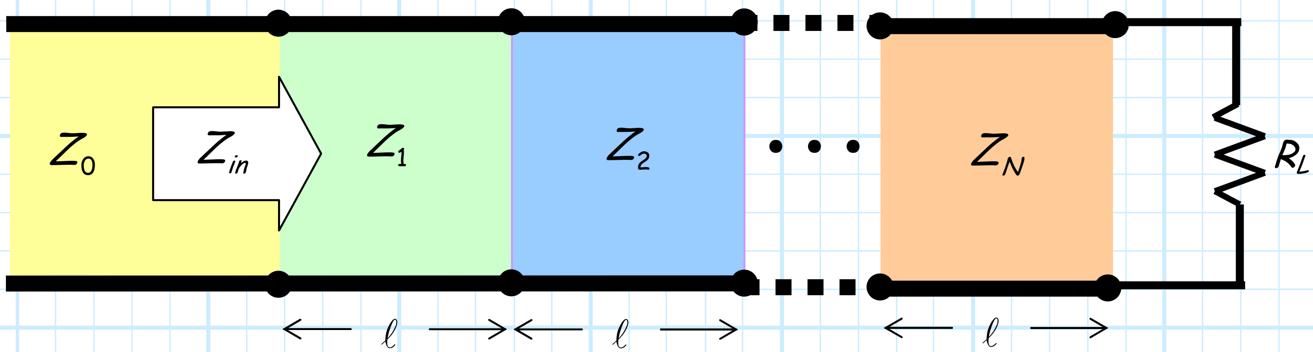
Set section lengths ℓ so that they are a **quarter-wavelength** ($\lambda_0/4$) at the design frequency ω_0.

And that pesky constant A ?

Q: *I see! And then we select all the values Z_n such that $\Gamma_n = A\, C_n^N$. But wait! **What** is the value of **A** ??*

A: *We can determine this value by evaluating a **boundary condition!***

Specifically, we can **easily** determine the value of $\Gamma(\omega)$ at $\omega = 0$.



Note as ω approaches **zero**, the electrical length $\beta\ell$ of each section will **likewise** approach zero. Thus, the input impedance Z_{in} will simply be equal to R_L as $\omega \rightarrow 0$.

As a result, the input reflection coefficient $\Gamma(\omega = 0)$ **must** be:

$$\Gamma(\omega = 0) = \frac{Z_{in}(\omega = 0) - Z_0}{Z_{in}(\omega = 0) + Z_0} = \frac{R_L - Z_0}{R_L + Z_0}$$

Aren't boundary conditions great ?

However, we **likewise** know that:

$$\begin{aligned}\Gamma(0) &= A\left(1+e^{-j2(0)}\right)^N \\ &= A(1+1)^N \\ &= A\,2^N\end{aligned}$$

Equating the two expressions:

$$\Gamma(0) = A\,2^N = \frac{R_L - Z_0}{R_L + Z_0}$$

And therefore:

$$A = 2^{-N}\frac{R_L - Z_0}{R_L + Z_0}$$

(A can be negative!)

We now have a form to calculate the **required marginal reflection coefficients** Γ_n:

$$\Gamma_n = A\,C_n^N = \frac{A\,N!}{(N-n)!\,n!}$$

How do I determine characteristic impedance?

Of course, we **also** know that these marginal reflection coefficients are physically related to the characteristic impedances of each section as:

$$\Gamma_n = \frac{Z_{n+1} - Z_n}{Z_{n+1} + Z_n}$$

Equating the two and solving, we find that that the section characteristic impedances must satisfy:

$$Z_{n+1} = Z_n \frac{1+\Gamma_n}{1-\Gamma_n} = Z_n \frac{1+A\,C_n^N}{1-A\,C_n^N}$$

Note this is an **iterative** result—we determine Z_1 from Z_0, Z_2 from Z_1, and so forth.

Q: *This result* ***appears*** *to be our second design equation. Is there some reason why you didn't draw a big blue box around it?*

A: Alas, there is a **big problem** with this result.

The BIG problem with this result!

Note that there are N+1 coefficients Γ_n (i.e., $n \in \{0,1,\cdots,N\}$) in the Binomial series, yet there are only N design degrees of freedom (i.e., there are only N transmission line sections!).

Thus, our design is a bit **over constrained**, a result that manifests itself the finally marginal reflection coefficient Γ_N.

Note from the iterative solution above, the **last** transmission line impedance Z_N is selected to satisfy the **mathematical** requirement of the **penultimate** reflection coefficient Γ_{N-1} :

$$\Gamma_{N-1} = \frac{Z_N - Z_{N-1}}{Z_N + Z_{N-1}} = A\, C_{N-1}^{N}$$

Thus the last impedance must be:

$$Z_N = Z_{N-1} \frac{1 + A\, C_{N-1}^{N}}{1 - A\, C_{N-1}^{N}}$$

Our degrees of freedom have run out!

But there is **one more** mathematical requirement! The last marginal reflection coefficient **must** likewise satisfy:

$$\Gamma_N = A\, C_N^N = 2^{-N} \frac{R_L - Z_0}{R_L + Z_0}$$

where we have used the fact that $C_N^N = 1$.

But, we **just** selected Z_N to satisfy the requirement for Γ_{N-1},—we have no **physical** design parameter to satisfy this last **mathematical** requirement!

As a result, we find to our great consternation that the last requirement is not satisfied:

$$\Gamma_N = \frac{R_L - Z_N}{R_L + Z_N} \neq A\, C_N^N \quad !!!!!!$$

Q: *Yikes! Does this mean that the resulting matching network will* ***not*** *have the desired Binomial frequency response?*

A: That's **exactly** what it means!

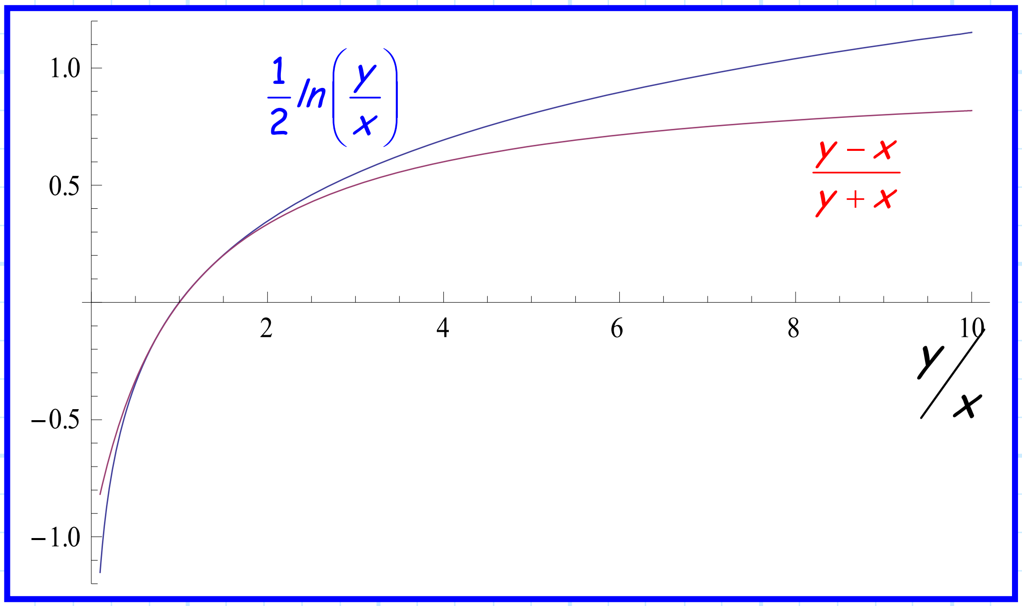
&*#@*!&!!!!!

Q: *You big #%@#$%&!!!! **Why** did you **waste** all my time by discussing an over-constrained design problem that can't be built?*

A: Relax; there is a **solution** to our dilemma—albeit an **approximate** one. **You** undoubtedly have previously used the **approximation**:

$$\frac{y-x}{y+x} \approx \frac{1}{2} ln\left(\frac{y}{x}\right)$$

An approximation that is especially **accurate** when $|y-x|$ is small (i.e., when $y/x \simeq 1$).



Use this approximation for value A!

Now, we know that the values of Z_{n+1} and Z_n in a multi-section matching network are typically **very close**, such that $\left|Z_{n+1} - Z_n\right|$ is small.

Thus, we use the approximation:

$$\Gamma_n = \frac{Z_{n+1} - Z_n}{Z_{n+1} + Z_n} \approx \frac{1}{2} ln\left(\frac{Z_{n+1}}{Z_n}\right)$$

Likewise, we can **also** apply this approximation (although not as accurately) to the value of A:

$$A = 2^{-N} \frac{R_L - Z_0}{R_L + Z_0} \approx 2^{-(N+1)} ln\left(\frac{R_L}{Z_0}\right)$$

Let's try this again—with approximations!

So, let's **start over**, only this time we'll use these **approximations**. First, determine A :

$$A \approx 2^{-(N+1)} \ln\left(\frac{R_L}{Z_0}\right)$$

(A can be negative!)

Now use this result to calculate the **mathematically required** marginal reflection coefficients Γ_n:

$$\Gamma_n = A\, C_n^N = \frac{A\; N!}{(N-n)!\,n!}$$

Here's (finally) our second design rule!

Of course, we **also** know that these marginal reflection coefficients are **physically** related to the **characteristic impedances** of each section as:

$$\Gamma_n \approx \frac{1}{2} ln\left(\frac{Z_{n+1}}{Z_n}\right)$$

Equating the two and solving, we find that that the section characteristic impedances **must** satisfy:

$$Z_{n+1} = Z_n \, exp\left[2\Gamma_n\right]$$

Now **this** is our **second design rule**. Note it is an **iterative** rule—we determine Z_1 from Z_0, Z_2 from Z_1, and so forth.

I don't understand what just happened

Q: *Huh? How is this any better? How does applying* ***approximate*** *math lead to a* ***better*** *design result??*

A: Applying these approximations help **resolve** our over-constrained problem. Recall that the over-constraint resulted in:

$$\Gamma_N = \frac{R_L - Z_N}{R_L + Z_N} \neq A\, C_N^N$$

But, as it turns out, these approximations leads to the **happy situation** where:

$$\Gamma_N \approx \frac{1}{2} ln\left(\frac{R_L}{Z_N}\right) = A\, C_N^N$$

← **A Sanity check!!**

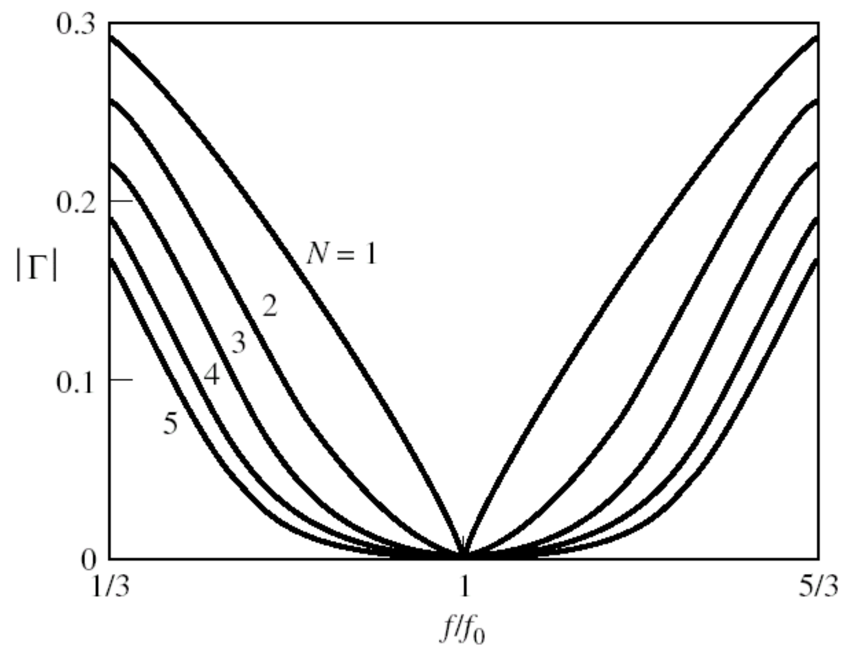
provided that the value A is likewise the **approximation** given above.

I *still* don't understand what just happened

Effectively, these approximations couple the results, such that each value of characteristic impedance Z_n **approximately** satisfies both Γ_n and Γ_{n+1}. Summarizing:

* If you use the **"exact"** design equations to determine the characteristic impedances Z_n, the **last** value Γ_N will exhibit a **significant** numeric error, and your design **will not** appear to be maximally flat.

* If you instead use the **"approximate"** design equations to determine the characteristic impedances Z_n, **all** values Γ_n will exhibit a **slight** error, but the resulting design **will** appear to be **maximally flat, Binomial** reflection coefficient function $\Gamma(\omega)$!



*Note that as we **increase** the number of **sections**, the matching **bandwidth** increases.*

Bandwidth: How do we define it?

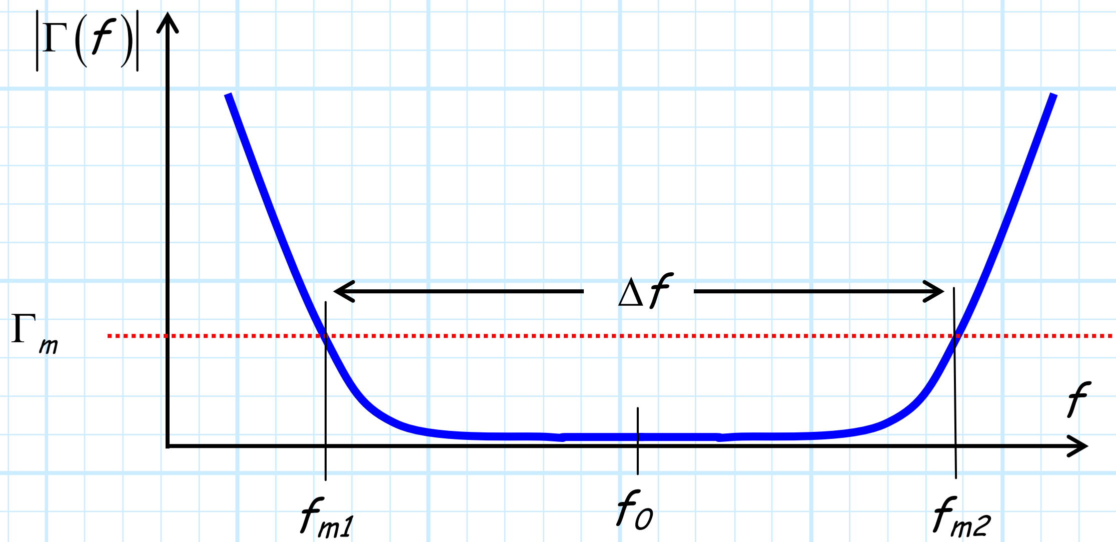
band·width (band′width′, -witth′) - noun

1. the range of frequencies within a ….

Q: *Can we determine the **value** of this bandwidth?*

A: Sure! But we first must **define** what we mean by bandwidth.

As we move from the design (perfect match) frequency f_0 the value $|\Gamma(f)|$ will **increase**. At some frequency (f_m, say) the magnitude of the reflection coefficient will increase to some **unacceptably** high value (Γ_m, say). At that point, we **no longer** consider the device to be matched.



Bandwidth: How do we calculate it?

Note there are **two** values of frequency f_m —one value **less** than design frequency f_0, and one value **greater** than design frequency f_0.

These two values define the **bandwidth** Δf of the matching network:

$$\Delta f = f_{m2} - f_{m1} = 2\left(f_0 - f_{m1}\right) = 2\left(f_{m2} - f_0\right)$$

Q: *So what is the* ***numerical*** *value of* Γ_m*?*

A: **I don't know**—it's up to **you** to decide!

Every engineer must determine what **they** consider to be an acceptable match (i.e., decide Γ_m).

This decision depends on the **application** involved, and the **specifications** of the overall microwave system being designed.

However, we **typically** set Γ_m to be 0.2 or less.

We get to perform some Algebra!!

Q: *OK, after we have selected Γ_m, can we determine the* ***two*** *frequencies f_m?*

A: Sure! We just have to do a little **algebra**.

We start by **rewriting** the Binomial function:

$$\begin{aligned}\Gamma(\theta) &= A\left(1+e^{-j2\theta}\right)^N \\ &= Ae^{-jN\theta}\left(e^{+j\theta}+e^{-j\theta}\right)^N \\ &= Ae^{-jN\theta}\left(e^{+j\theta}+e^{-j\theta}\right)^N \\ &= Ae^{-jN\theta}\left(2\cos\theta\right)^N\end{aligned}$$

Now, we take the **magnitude** of this function:

$$\begin{aligned}\left|\Gamma(\theta)\right| &= 2^N\left|A\right|\left|e^{-jN\theta}\right|\left|\cos\theta\right|^N \\ &= 2^N\left|A\right|\left|\cos\theta\right|^N\end{aligned}$$

It gets better—even more algebra!!

Now, we **define** the values θ where $|\Gamma(\theta)| = \Gamma_m$ as θ_m. I.E., :

$$\begin{aligned} \Gamma_m &= |\Gamma(\theta = \theta_m)| \\ &= 2^N |A| |\cos\theta_m|^N \end{aligned}$$

We can now solve for θ_m (in **radians**!) in terms of Γ_m:

$$\theta_{m1} = \cos^{-1}\left[\frac{1}{2}\left(\frac{\Gamma_m}{|A|}\right)^{1/N}\right] \qquad \theta_{m2} = \cos^{-1}\left[-\frac{1}{2}\left(\frac{\Gamma_m}{|A|}\right)^{1/N}\right]$$

Note that there are **two solutions** to the above equation (one **less** that $\pi/2$ and one **greater** than $\pi/2$)!

Now, we can **convert** the values of θ_m into specific **frequencies**.

Converting θ_m to f_m

Recall that $\omega T = \theta$, therefore:

$$\omega_m = \frac{1}{T}\theta_m = \frac{v_p}{\ell}\theta_m$$

But recall also that $\ell = \lambda_0/4$, where λ_0 is the wavelength at the **design frequency** f_0 (not f_m!), and where $\lambda_0 = v_p/f_0$.

Thus we can conclude:

$$\omega_m = \frac{v_p}{\ell}\theta_m = \frac{4v_p}{\lambda_0}\theta_m = (4f_0)\theta_m$$

or:

$$f_m = \frac{1}{2\pi}\frac{v_p}{\ell}\theta_m = \frac{(4f_0)\theta_m}{2\pi} = \frac{(2f_0)\theta_m}{\pi}$$

where θ_m is expressed in **radians**.

And thus the bandwidth is…

Therefore:

$$f_{m1} = \frac{2f_0}{\pi} \cos^{-1}\left[+\frac{1}{2}\left(\frac{\Gamma_m}{|A|}\right)^{1/N} \right] \qquad f_{m2} = \frac{2f_0}{\pi} \cos^{-1}\left[-\frac{1}{2}\left(\frac{\Gamma_m}{|A|}\right)^{1/N} \right]$$

Thus, the **bandwidth** of the binomial matching network can be determined as:

$$\begin{aligned} \Delta f &= 2\left(f_0 - f_{m1}\right) \\ &= 2f_0 - \frac{4f_0}{\pi} \cos^{-1}\left[+\frac{1}{2}\left(\frac{\Gamma_m}{|A|}\right)^{1/N} \right] \end{aligned}$$

Note that this equation can be used to determine the **bandwidth** of a binomial matching network, given Γ_m and number of sections N.

However, it can likewise be used to determine the **number of sections** N required to meet a specific bandwidth requirement!

In summary, our design steps

Finally, we can list the **design steps** for a binomial matching network:

1. **Determine** the value N required to meet the bandwidth (Δf and Γ_m) requirements.

2. Determine the **approximate** value A from Z_0, R_L and N.

3. Determine the **marginal reflection coefficients** $\Gamma_n = A\,C_n^N$ required by the **binomial** function.

4. Determine the characteristic impedance of each section using the **iterative approximation**:

$$Z_{n+1} = Z_n \, exp\left[2\Gamma_n\right]$$

5. Perform the **sanity check**:

$$\Gamma_N \approx \frac{1}{2} ln\left(\frac{R_L}{Z_N}\right) = A\,C_N^N$$

6. Determine section **length** $\ell = \lambda_0/4$ for design frequency f_0.