# **Interpreter Optimizations**

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# Outline

- Background
- Context threading
- Superoperators
- Conclusion

#### Background

- Decode Dispatch Interpreter high interpretation cost
- Indirect Threaded Interpretation dispatch code to interpreter routines
- Direct Threaded Interpretation
  - Best choice amongst three

# Background

Motivation:

- Direct-threaded Indirect branches to dispatch bytecodes
- Deeply-pipelined architecture rely on branch predictors for performance
- context problem indirect branches are poorly predicted

Idea:

- Refine the dispatch itself **context threading**
- Efficient or fewer dispatches **superoperators**

# **Context Threading**

Idea:

- Align hardware and virtual machine state
- i.e. correlate native pc (tpc) to vpc (spc)
- Improve branch prediction

#### Context Threading - direct threaded interpreter



#### Direct threaded interpreter



#### Diagram credits:

http://www.ittc.ku.edu/~kulkarni/teaching/EECS768/slides/chapter2.pdf

Diagram credits: http://www.cs.toronto.edu/~matz/pubs/demkea\_context.pdf paper

#### Context Threading - direct threaded interpreter

But there is a problem with this approach

- Target of the dispatch branch depends on vpc, but not hardware pc
- context problem

#### Context threading - direct threaded interpreter



Image credits: <u>http://www.cs.toronto.edu/syslab/talks/cgoAsWorkshop.pdf</u>

#### Context threading - direct threaded interpreter

But there is a problem with this approach

- Target of the dispatch branch depends on vpc, but not hardware pc
- context problem

#### Leverage hardware predictors

## Context threading

Virtual program may contain the following control flow types:

- Conditional branches
- Call and returns
- Unconditional branches
- Linear code

But direct threading uses indirect branches for all types of control flows

Idea: expose the virtual control flow patterns to the hardware, map spc and tpc

# Context threading

- Handling linear dispatch
- Handling virtual branches
- handling virtual call and return

#### Context threading - Handling the control flow



 Handling linear dispatch

Diagram credits: <u>http://www.cs.toronto.edu/~matz/pubs/demkea\_context.pdf</u>

#### Context threading - Handling the control flow



Handling virtual branches

• Inline the bodies into CTT

Handling virtual call and return

 new handler instructions are added

Diagram credits: <u>http://www.cs.toronto.edu/~matz/pubs/demkea\_context.pdf</u>

# Context threading - Tiny inlining

- Context threading can be combined with inlining strategies
- Inline techniques can be used to reduce the pipeline hazards
- simple heuristic
- Inlining all the small bodies
- This removes dispatch overhead surrounding smallest bodies
- Further optimizing the context threaded interpreter

### **Context threading - Results**



95% branch mispredictions eliminated on average

27% reduction in execution time on average

Results: https://webdocs.cs.ualberta.ca/~amaral/cascon/CDP05/slides/CDP05-berndl.pdf

- Efficient or fewer dispatches as seen earlier
- An optimization technique for bytecode interpreters
- combines smaller atomic operations
- Decrease the executable size
- Increase speed

- simple e.g.:
- ADD(R, 4)
- Requires several virtual instructions
- Extend the virtual instruction set to add a single "super-instruction"
- Decreases the executable size and makes the interpreter faster
- Icc ir uses 109 operators
- why not use remaining 147 byte codes for superoperators

Inference Heuristic:

- Interpreter includes a heuristic method for inferring a good set of superoperators
- Reads IR trees and decides which adjacent operators to merge
- Each tree is weighted to guide the heuristic
- Optimize for space and time

Let's take an example.

- 1. Assume inputs trees with weights.
- 2. Now find the operator frequencies.
- 3. Now find the frequency of adjacent nodes.
- 4. Find the node with the highest frequency.
- 5. Repeat the operation.

I(A(Z,Y))	10	I(B(*))	10
Α(Υ,Υ)	1	B(Z)	10
Y	12	B(Y)	1
Z	10		20.20
I	10	I(C)	10
A	11	B(Y)	1
I(A(*))	10	Y	1
A(Z,*)	10	Z	0
A(*,Y)	11	I	10
A(Y,*)	1	A	0
I(B(Z))	10	В	1
B(Y)	1	С	10

#### Conclusion

- Context problem: Branch mispredictions due to the mismatch between virtual and native control flow
- Context threading method
- Superoperators optimize the interpreters by reducing branches

#### Thanks for listening.